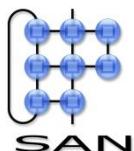


Generalized fixed-priority scheduling with limited preemptions

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Where innovation starts



Motivation

- **Drawbacks of arbitrary preemptions (FPPS):**
 - worst-case memory requirements;
 - cost of context switching (e.g. cache).
- **Non-preemptive scheduling (FPNS):**
 - resolved by FPNS,
 - at the cost of lower schedulability.
- **Alternative refinements of FPPS (FPTS & FPDS):**
 - reduced memory costs (compared to FPPS);
 - improved schedulability (compared to FPPS);
 - orthogonal approaches !

Motivation

- **Goal:**
 - combine strength of FPTS and FPDS in a single scheme: **FPGS**,
 - with the aim to improve efficiency,
 - focus on improvement of the feasibility.

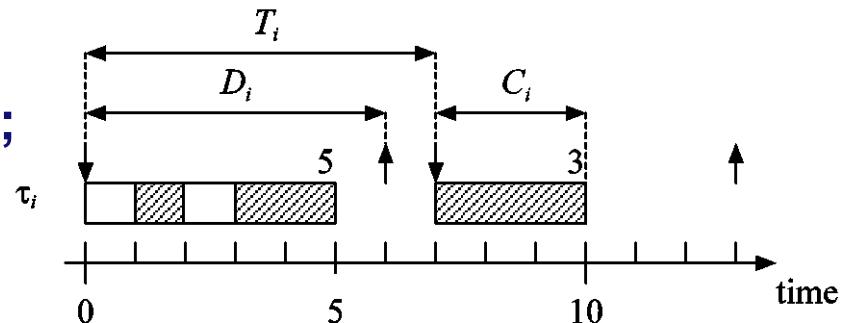
Overview

- Motivation
- **Introduction fixed-priority scheduling (FPS)***
 - model
 - FPPS and FPNS
- **Existing limited-preemptive algorithms***
 - preemption thresholds (FPTS)
 - deferred preemption (FPDS)
- **Novel hybrid algorithms**
- **Conclusion**

* Buttazzo, Bertogna, and Yao, IEEE TII, 2012.

Introduction FPS – model

- Events: implicit
- Tasks (τ):
 - independent, no self-suspension
 - characteristics (R^+):
 - minimal inter-arrival time (T);
 - computation time (C);
 - deadline (D);
- Scheduling algorithm:
 - fixed-priority (π) & non-idling;
 - [non-] preemptive
- Platform: single CPU



↓ activation



interference or blocking

↑ deadline



execution

Introduction FPS – FPPS and FPNS

- **FPPS:**
 - highest priority task with work pending executes;
 - a task experiences *interference* from *higher* priority tasks (due to *delays* and *preemptions*).
- **FPNS:**
 - tasks run to completion;
 - the highest priority task is selected to run next;
 - a task experiences:
 - *interference* from *higher* priority tasks (*only delays*);
 - *blocking* from *lower* priority tasks.

Introduction FPS – FPPS and FPNS

- **FPNS:**
 - **blocking:** $B_i = \max(0, \max_{l:\pi_i > \pi_l} C_l)$
- ***Blocking tolerance* (β_i) [1]:**
 - the maximum amount of time that a task (τ_i) can be blocked without missing its deadline (D_i);
 - depends on scheduling algorithm.
- **Neither FPPS dominates FPNS nor vice versa.**

[1] V.B. Lortz and K.G. Shin, IEEE TSE, 1995.

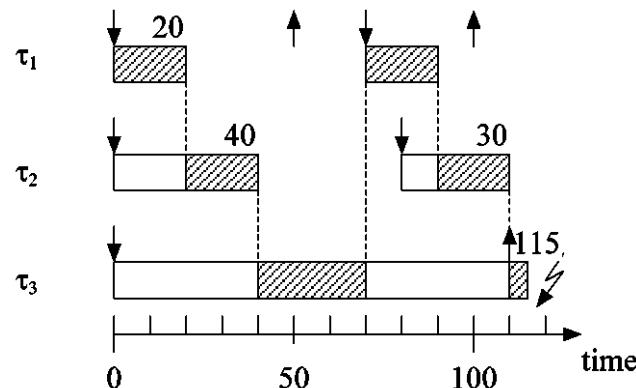
Existing limited-preemptive algorithms

- Two orthogonal approaches: FPTS [2, 3] and FPDS
 - FPTS: *preemption threshold* ($\theta_i \geq \pi_i$)
 - interference: (reduce *preemptions*)
 - tasks τ_h with $\pi_h > \theta_i$ can preempt τ_i ;
 - blocking: $B_i = \max(0, \max_{l: \theta_l \geq \pi_i > \pi_l} C_l)$
 - Special cases of FPTS:
 - FPPS: $\theta_i = \pi_i$;
 - FPNS: $\theta_i = \pi_1$.

- [2] Y. Wang and M. Saksena, RTCSA, 1999.
[3] J. Regehr, RTSS, 2002.

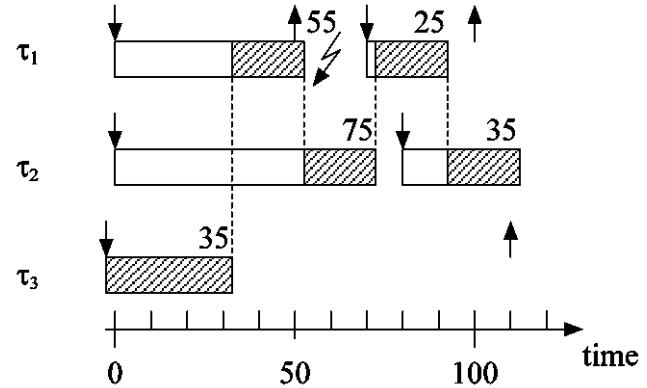
FPTS – preemption thresholds

	T_i	D_i	C_i	π_i	WR_i^P	WR_i^N
τ_1	70	50	20	3	20	55
τ_2	80	80	20	2	40	75
τ_3	200	100	35	1	115	75



Not schedulable with FPPS

- Blocking tolerance τ_1 :
- $\beta_1 = 30, C_2 < \beta_1 < C_3$.
- Blocking tolerance τ_2 :
- FPPS: $\beta_2^P = 30 < C_3$;
 - FPNS: $\beta_2^N = 40 > C_3$.



Not schedulable with FPNS

FPTS – preemption thresholds

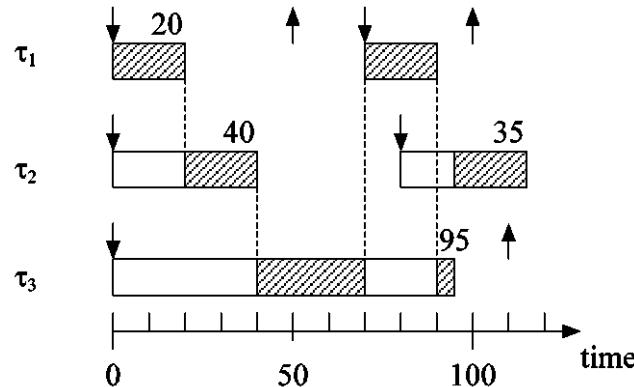
	T_i	D_i	C_i	π_i	WR_i^P	WR_i^N	θ_i	WR_i^T
τ_1	70	50	20	3	20	55	3	40
τ_2	80	80	20	2	40	75	3	75
τ_3	200	100	35	1	115	75	2	95

Blocking tolerance τ_1 :

- $\beta_1 = 30, C_2 < \beta_1 < C_3$.

Blocking tolerance τ_2 :

- FPPS: $\beta_2^P = 30 < C_3$;
- FPNS: $\beta_2^N = 40 > C_3$.



Schedulable with FPTS

Existing limited-preemptive algorithms

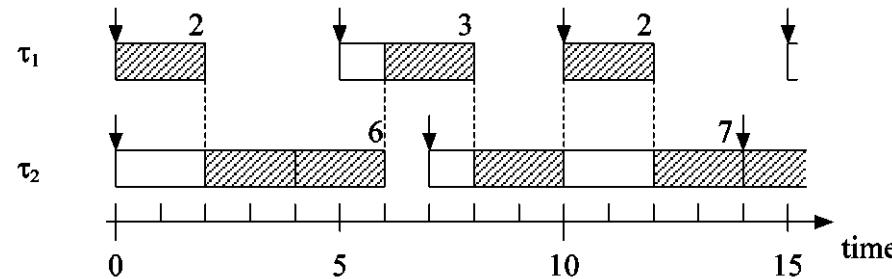
- Two orthogonal approaches: FPTS and FPDS [4, 5]
 - FPDS: *deferred preemption*
 - a task is a sequence of m_i non-preemptive *sub-tasks*;
 - characteristic subtask $\tau_{i,k}$: computation time $C_{i,k}$;
 - tasks τ_h with $\pi_h > \pi_i$ can only preempt τ_i at *preemption points* (between subtasks);
 - blocking: $B_i = \max(0, \max_{l:\pi_l > \pi_i} \max_{1 \leq k \leq m_l} C_{l,k})$
 - Special case of FPDS:
 - FPNS: $m_i = 1$.

- [4] A. Burns, in *Advances in Real-Time Systems*, 1994
[5] R. Bril, J. Lukkien, and W. Verhaegh, ECRTS, 2007.

FPDS – deferred preemption

	$T_i = D_i$	C_i	π_i	WR_i^P	WR_i^N	WR_i^D
τ_1	5	2	2	2	6	4
τ_2	7	2+2	1	8	6	7

Blocking tolerance τ_1 :
• $\beta_1 = 3 < C_3$.



Schedulable with FPDS

FPDS – deferred preemption

	$T_i = D_i$	C_i	π_i	WR_i^P	WR_i^N	WR_i^D
τ_1	5	2	2	2	6	4
τ_2	7	2+2	1	8	6	7

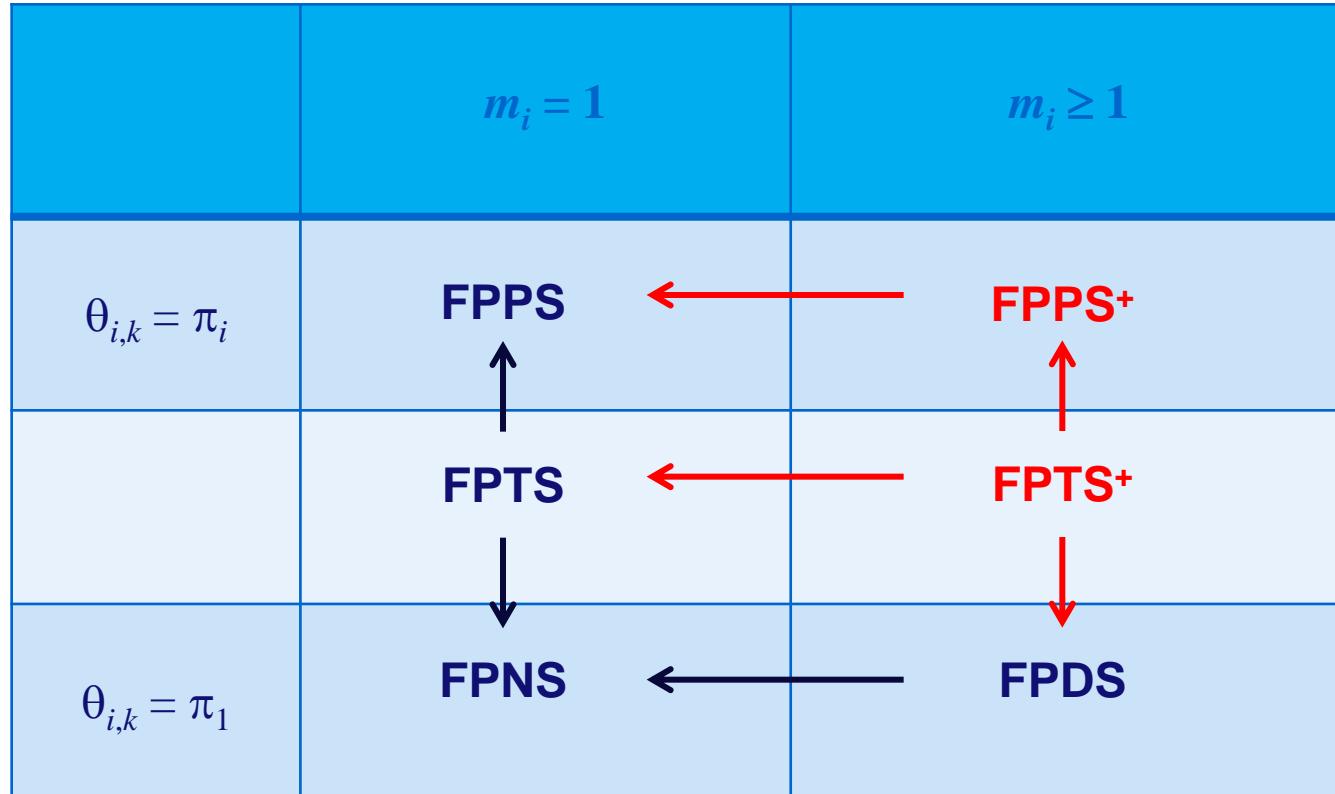
Blocking tolerance τ_1 :
• $\beta_1 = 3 < C_3$.

- **FPTS:**
 - $\theta_2 = \pi_2$: FPPS
 - $\theta_2 = \pi_1$: FPNS
- **Conclusion:**
 - **Not schedulable with FPTS, hence**
 - **FPTS does not dominate FPDS**

FPDS does not dominate FPTS

- Previous slide:
 - FPTS does not dominate FPDS;
- Without example (space & time reasons):
 - FPDS does not dominate FPTS
- Conclusion:
 - neither FPDS dominates FPTS nor vice versa

Novel hybrid algorithms



Generalization graph for FPS algorithms

Novel hybrid algorithms – model

- **FPTS⁺ [6, 7]**
 - a task is a sequence of m_i sub-tasks;
 - each subtask $\tau_{i,k}$ has a preemption threshold $\theta_{i,k}$;
 - a task has no (longer a) preemption threshold;
 - **blocking:** $B_i = \max(0, \max_{l: \pi_i > \pi_l} \max_{1 \leq k \leq m_l : \theta_{l,k} \geq \pi_i} C_{l,k})$
- **Special cases for FPPS⁺:**
 - $m_i = 1$: **FPTS**;
 - $\theta_{i,k} = \pi_i$: **FPPS⁺**;
 - $\theta_{i,k} = \pi_1$: **FPDS**.

[6] U. Keskin, R.J. Bril, J.J. Lukkien, ETFA, 2010.

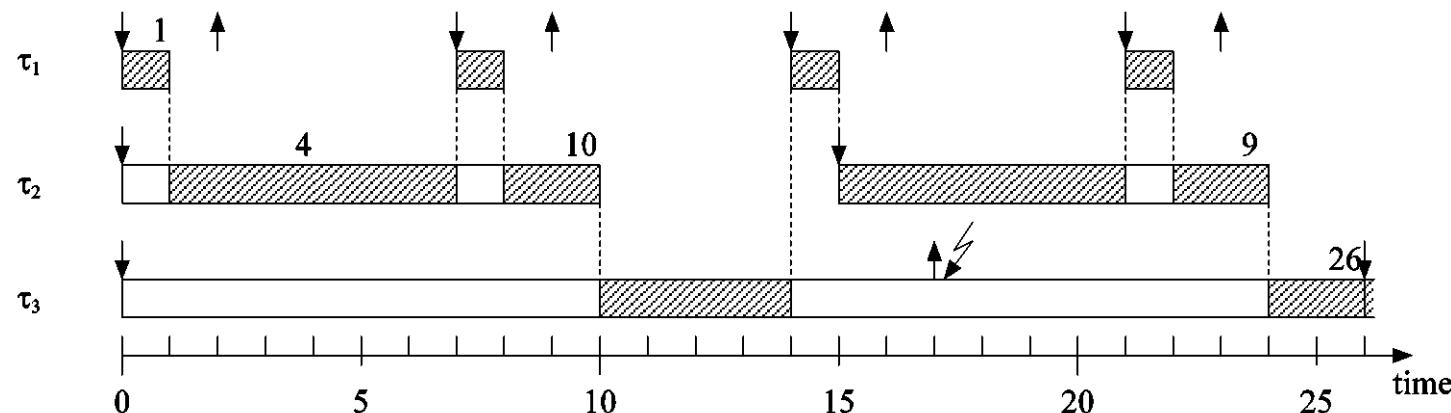
[7] G. Yao and G. Buttazzo, EMSOFT, 2010.

FPTS⁺ – preemption thresholds

	T_i	D_i	C_i	π_i	WR_i^P	WR_i^N
τ_1	7	2	1	3	1	9
τ_2	15	15	7+1	2	10	15
τ_3	26	17	1+5	1	26	16

Blocking tolerance τ_1 :

- $\beta_1 = 1, C_3 > \beta_1, C_2 > \beta_1.$



Blocking tolerance $\beta_1 = 1 \Rightarrow$ **Not** schedulable with FPDS.

FPTS⁺ – preemption thresholds

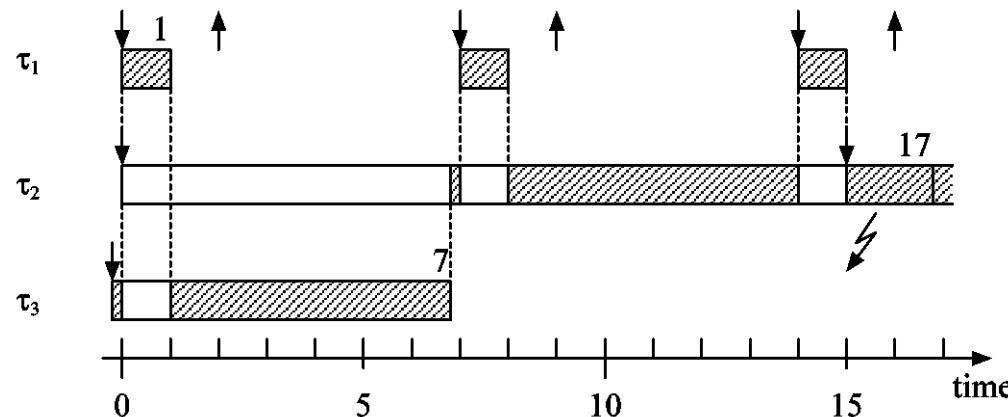
	T_i	D_i	C_i	π_i	WR_i^P	WR_i^N
τ_1	7	2	1	3	1	9
τ_2	15	15	7+1	2	10	15
τ_3	26	17	1+5	1	26	16

Blocking tolerance τ_1 :

- $\beta_1 = 1, C_3 > \beta_1, C_2 > \beta_1.$

Blocking tolerance τ_2 :

- FPPS: $\beta_2^P = 4 < C_3;$



Blocking tolerance $\beta_2 < C_3 \Rightarrow$ **Not** schedulable with FPTS.

FPTS⁺ – preemption thresholds

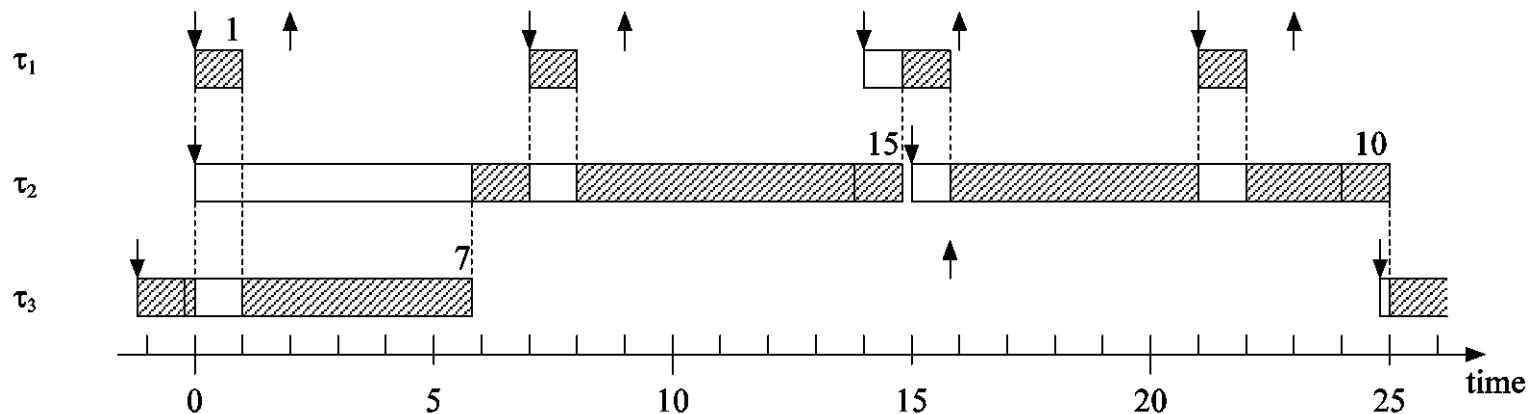
	T_i	D_i	C_i	π_i	WR_i^P	WR_i^N	$\theta_{i,1}$	$\theta_{i,2}$	WR_i^{T+}
τ_1	7	2	1	3	1	9	3	-	2
τ_2	15	15	7+1	2	10	15	2	3	15
τ_3	26	17	1+5	1	26	16	1	2	17

Blocking tolerance τ_1 :

- $\beta_1 = 1, C_3 > \beta_1, C_2 > \beta_1.$

Blocking tolerance τ_2 :

- FPPS: $\beta_2^P = 4 < C_3;$
- FPTS⁺: $\beta_2^{T+} = 5$



FPTS⁺ – preemption thresholds

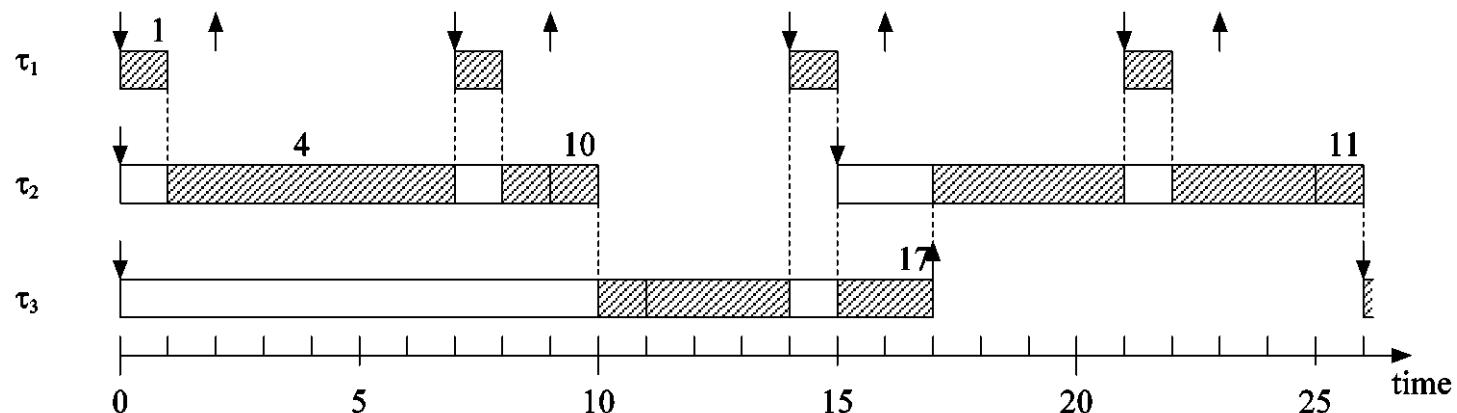
	T_i	D_i	C_i	π_i	WR_i^P	WR_i^N	$\theta_{i,1}$	$\theta_{i,2}$	WR_i^{T+}
τ_1	7	2	1	3	1	9	3	-	2
τ_2	15	15	7+1	2	10	15	2	3	15
τ_3	26	17	1+5	1	26	16	1	2	17

Blocking tolerance τ_1 :

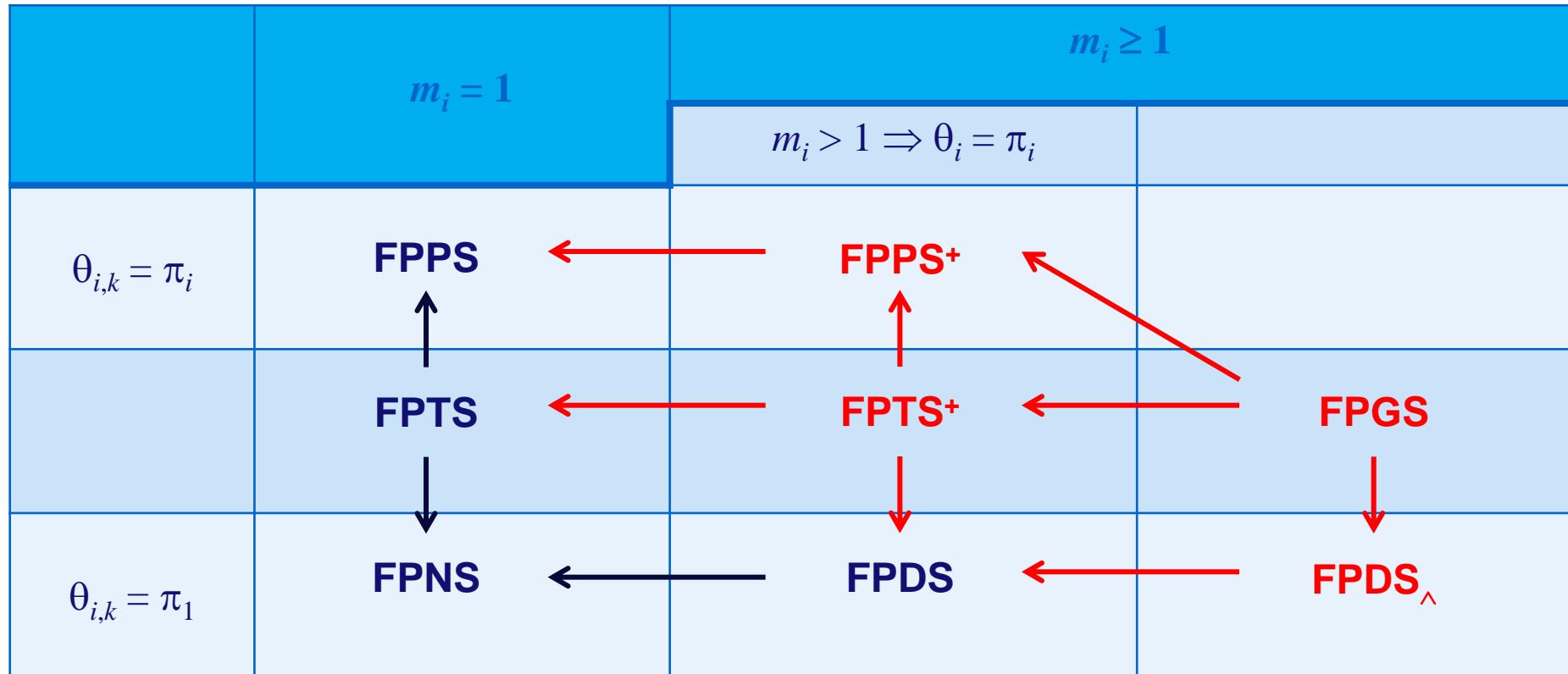
- $\beta_1 = 1, C_3 > \beta_1, C_2 > \beta_1.$

Blocking tolerance τ_2 :

- FPPS: $\beta_2^P = 4 < C_3;$
- FPTS⁺: $\beta_2^{T+} = 5$



Novel hybrid algorithms



Generalization graph for FPS algorithms

Novel hybrid algorithms – results

1. Novel scheduling algorithms: FPGS

- subtasks (similar to FPDS);
- preemption thresholds for tasks (FPTS) and subtasks;
- generalizes existing FPS algs.

2. Schedulability analysis for FPGS

- specializes to all existing FPS algs;

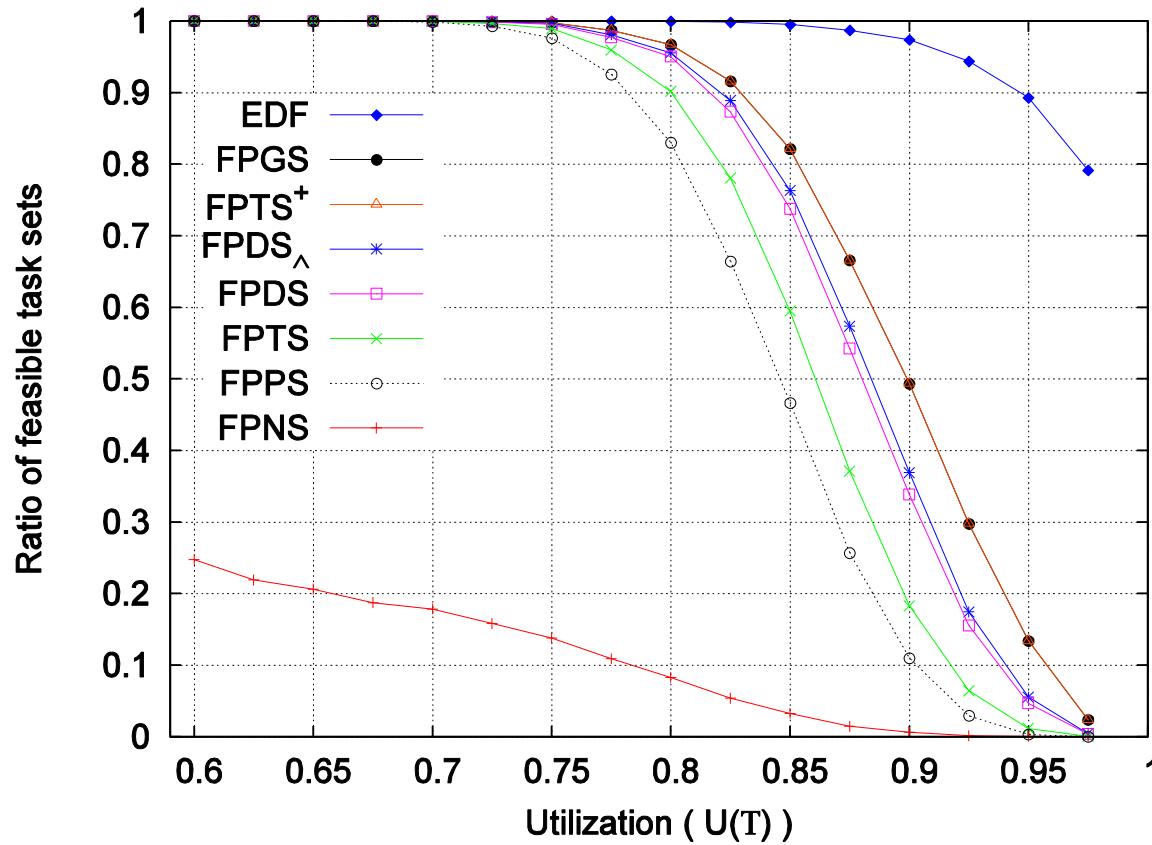
3. Algorithm to maximize schedulability under FPS:

- given: T_i , D_i , C_i , and π_i ;
- determine: $C_{i,mi}$, θ_i , $\theta_{i,mi}$ (inspired by [8]).

4. Evaluation: FPGS dominates existing FPS algs.

[8] M. Bertogna, G.C. Buttazzo, G, Yao, RTSS, 2011.

Novel hybrid algorithms – evaluation



10 tasks, 10.000 task sets,
 $T_i \in [100, 10.000]$ (uniform), U_i by UUnifast ($\Rightarrow C_i$),
 $D_i \in [0.5(T_i + C_i), T_i]$ (uniform);

Conclusion

- **FPGS and existing FPS algorithms:**
 - FPGS generalizes all others;
 - analysis of FPGS specializes to all others;
 - FPGS dominates all others.
- **Future work:**
 - further improvements of schedulability:
 - preemption thresholds for *preemption points*;
 - context switching cost;
 - ...