

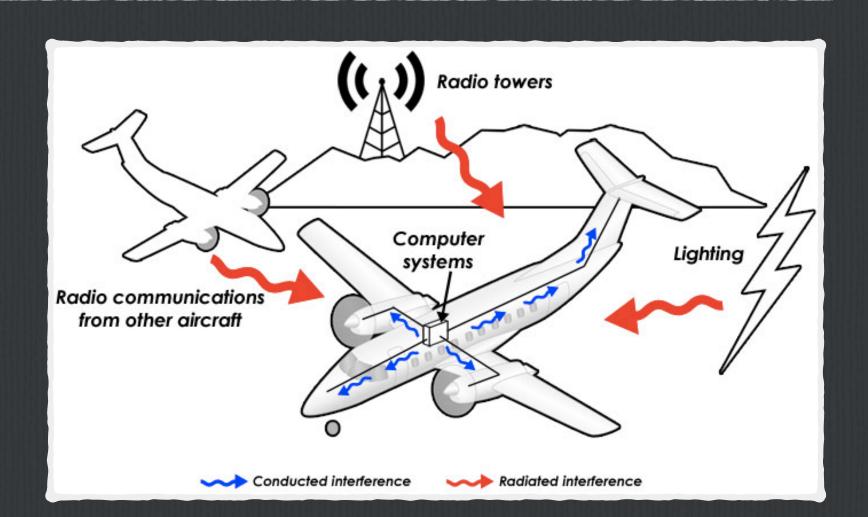
Quantifying the Resiliency of Fail-Operational Real-Time Networked Control Systems

Arpan Gujarati, Mitra Nasri, Björn B. Brandenburg



Embedded systems are susceptible to environmentally-induced transient faults

- ☐ Harsh environments
 - **→** Robots operating under hard radiation
 - **→** Industrial systems near high-power machinery
 - **➡ Electric motors** inside automobile systems
- ☐ Bit-flips in registers, buffers, network



Example*

*Mancuso R. Next-generation safety-critical systems on multi-core platforms. PhD thesis, UIUC, 2017.

- **→** One bit-flip in a 1 MB SRAM every 10¹² hours of operation
- → 0.5 billion cars with an average daily operation time of 5%
- **→** About 5,000 cars are affected by a bit-flip every day

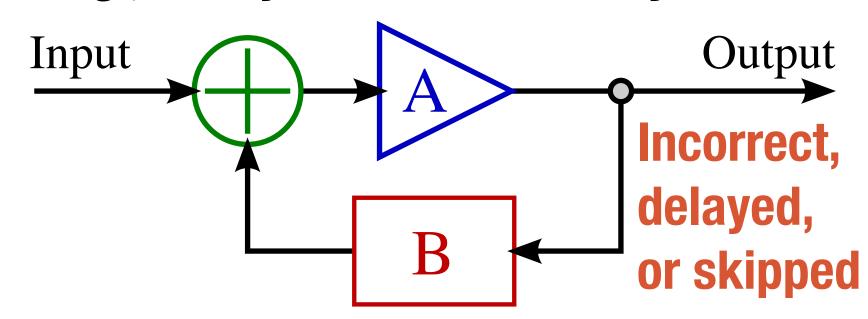
Failures and errors due to transient faults in distributed real-time systems

- ☐ Transmission errors
 - **→** Faults on the network
- ☐ Omission Errors
 - **→** Fault-induced kernel panics
- ☐ Incorrect computation Errors
 - **→** Faults in the memory buffers

Failures in:

- value domain (incorrect outputs)
- time domain (deadline violations)

E.g., safety-critical control system



Mitigating the effects of transient faults in distributed real-time systems

- ☐ Transmission errors
 - **→** Faults on the network
- Omission Errors
 - **→** Fault-induced kernel panics
- ☐ Incorrect computation Errors
 - **→** Faults in the memory buffers

Retransmissions at the network layer

Dual modular redundancy (DMR)

Triple modular redundancy (TMR)

Mitigating the effects of transient faults in distributed real-time systems

How can we objectively compare the reliability offered by different mitigation techniques?

- □ Omission Errors
 - **→** Fault-induced kernel panics
- ☐ Incorrect computation Errors
 - **→** Faults in the memory buffers

Retransmissions at the network layer

Dual modular redundancy (DMR)

Triple modular redundancy (TMR)

Mitigating the effects of transient faults in distributed real-time systems

- How does the real-time requirement affect system reliability?

 When does it really become a bottleneck?
- ☐ Omission Errors
 - **→** Fault-induced kernel panics

Dual modular redundancy (DMR)

What if the system is weakly-hard real-time, i.e., it can tolerate a few failures?

ular redundancy (TMR)

Problem: Reliability analysis of networked control systems

Given

- 1) Networked control system (messages, period)
- (2) Robustness specification (weakly-hard constraints)
- (3) Active replication scheme (DMR, TMR, others)
- 4) Peak transient fault rates (for the network and the hosts)

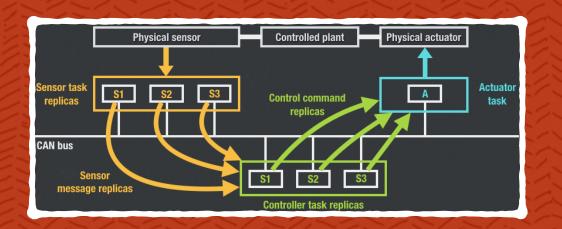
Objective

A safe upper bound on the failure rate of the networked control system \

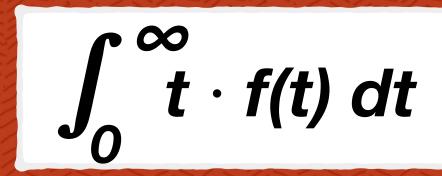
Failures-In-Time (FIT) = Expected # failures in one billion operating hours

Outline

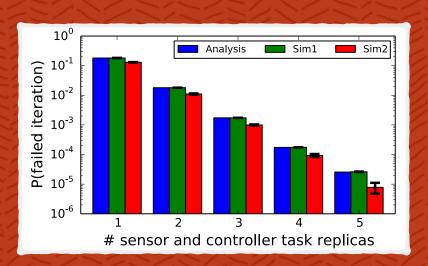
Analysis of a Controller Area Network (CAN) based networked control system



System Model



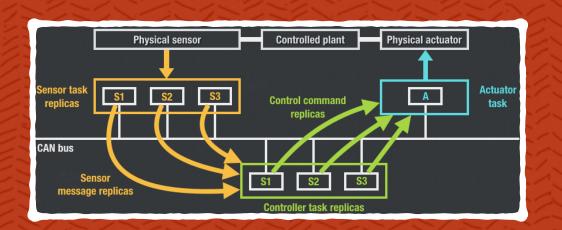
Analysis



Evaluation

Outline

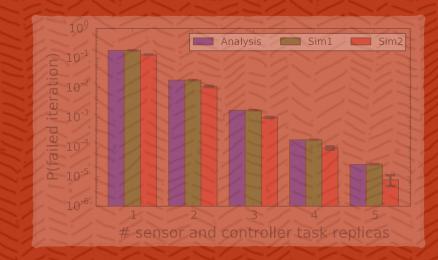
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System Model

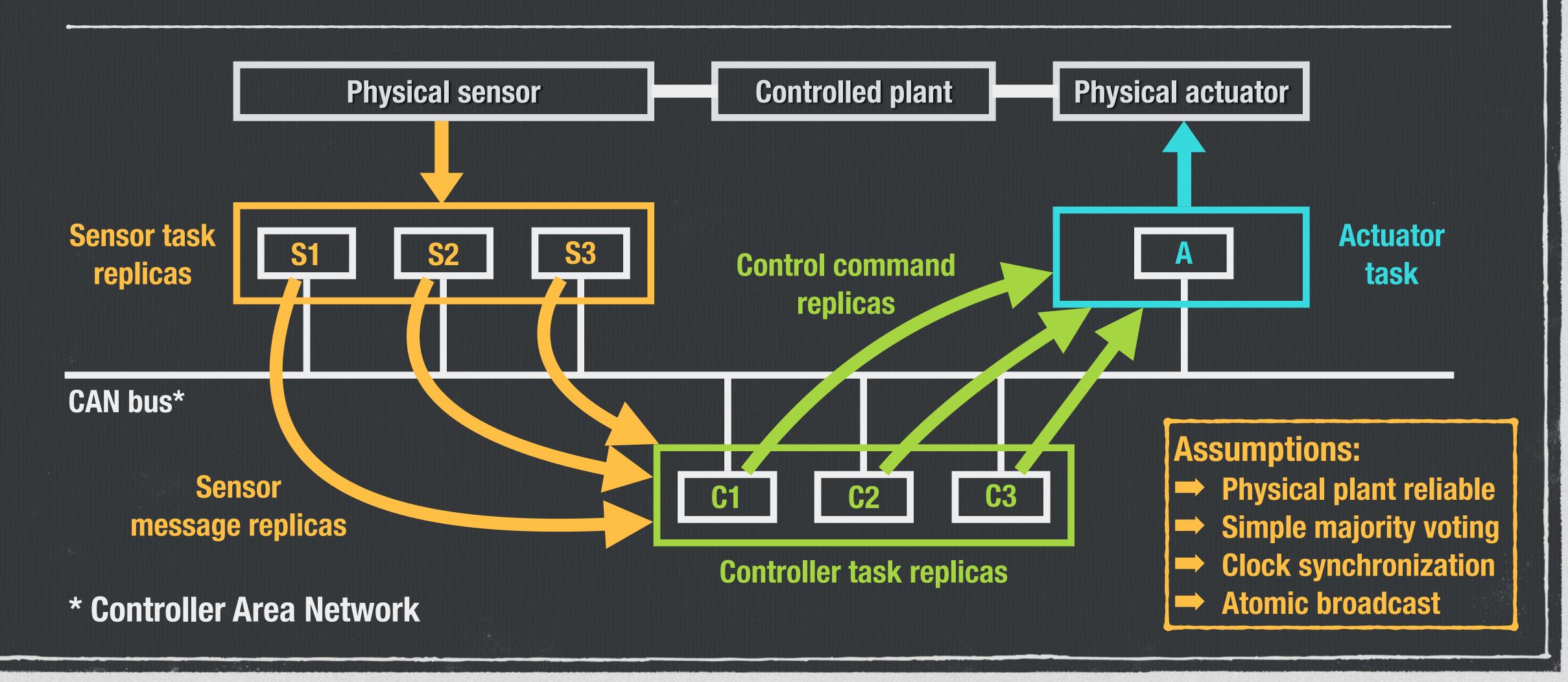


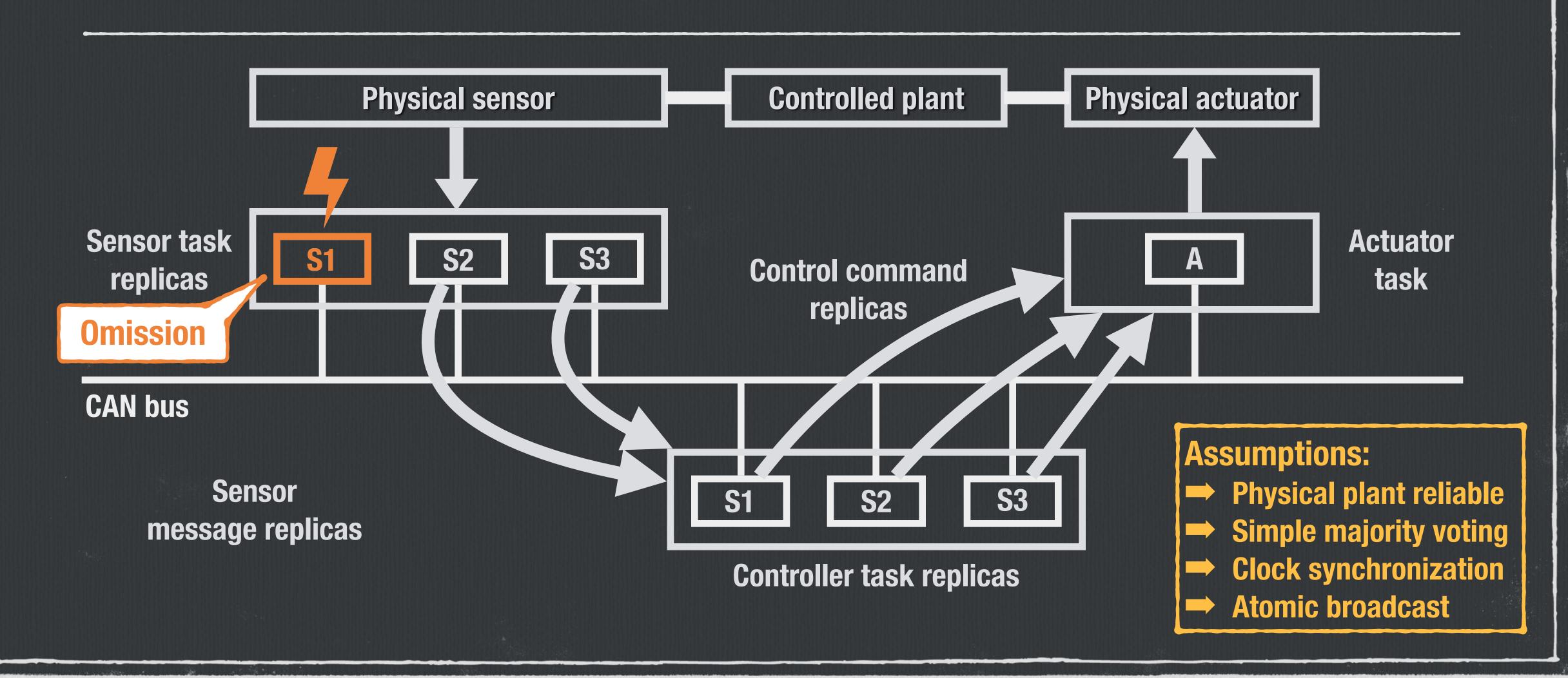
Analysis

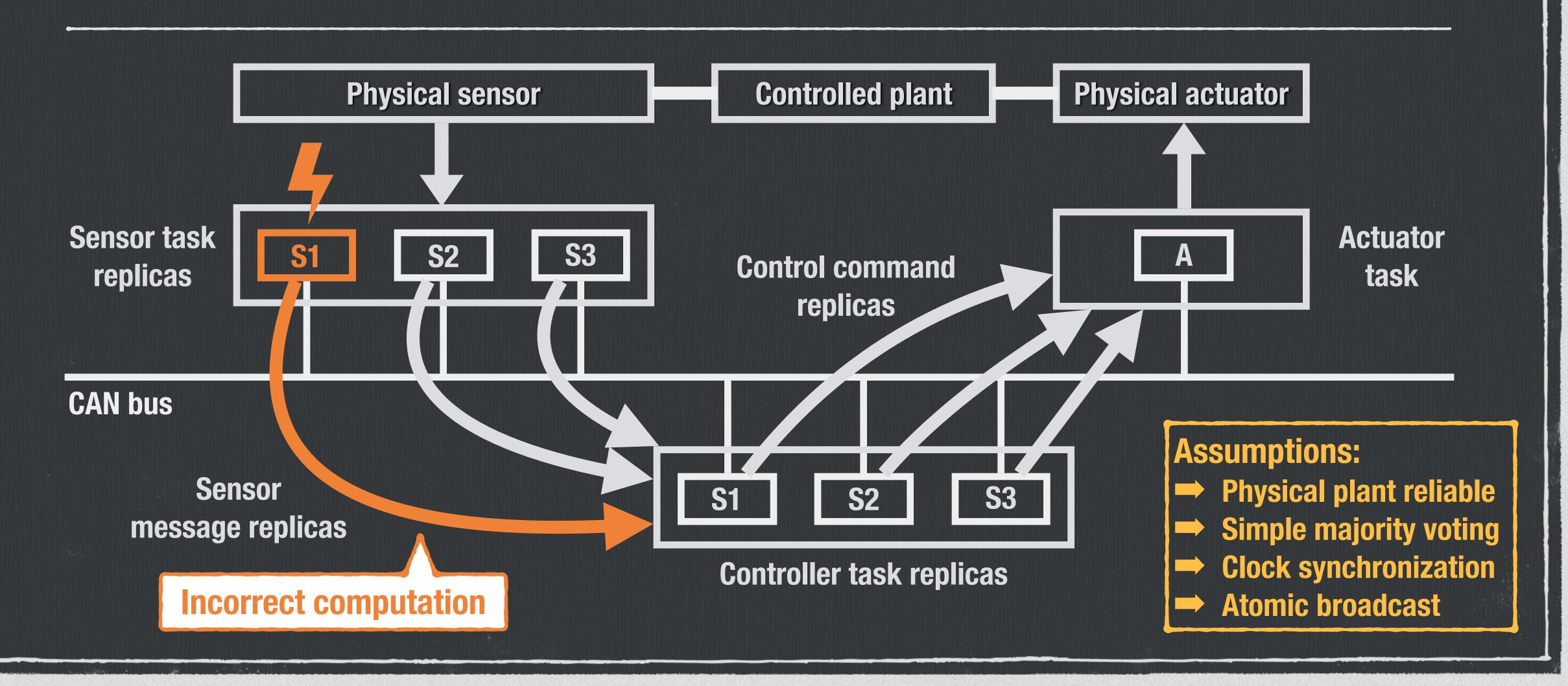


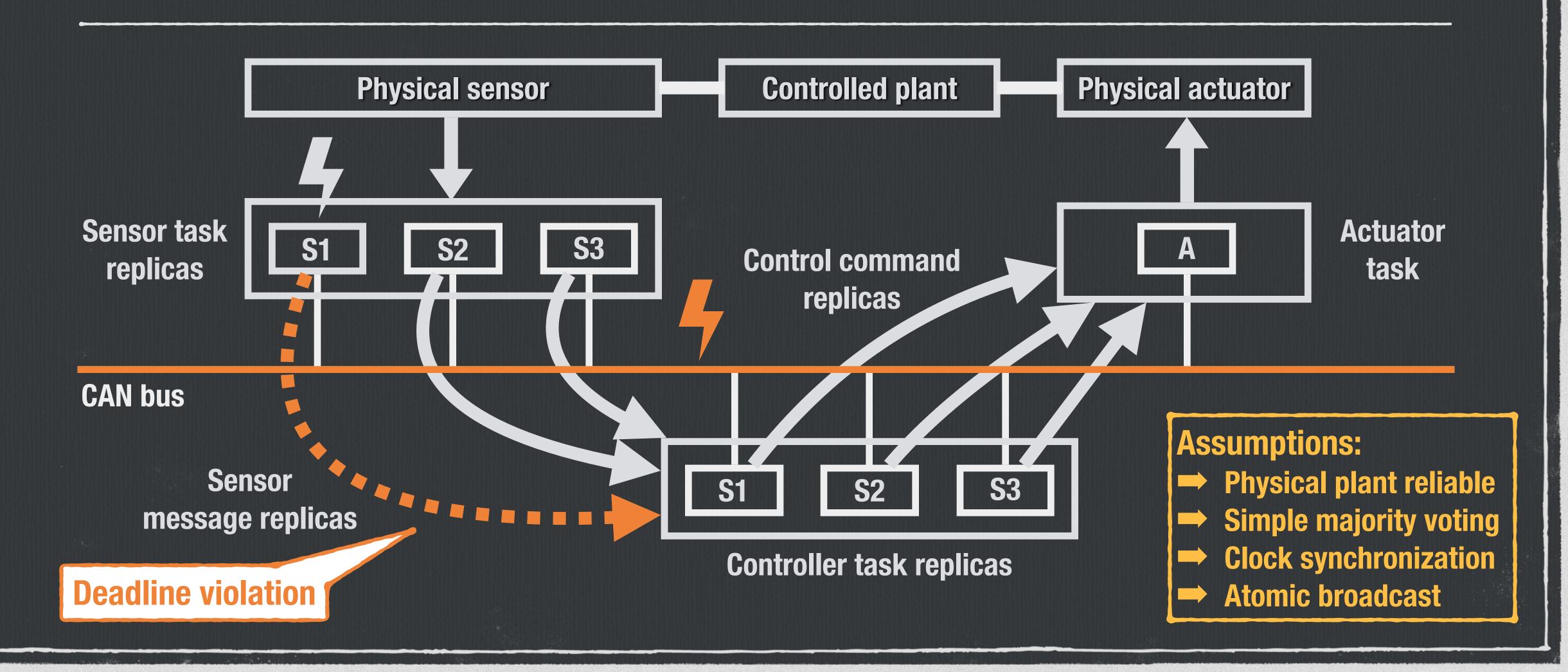
Evaluation

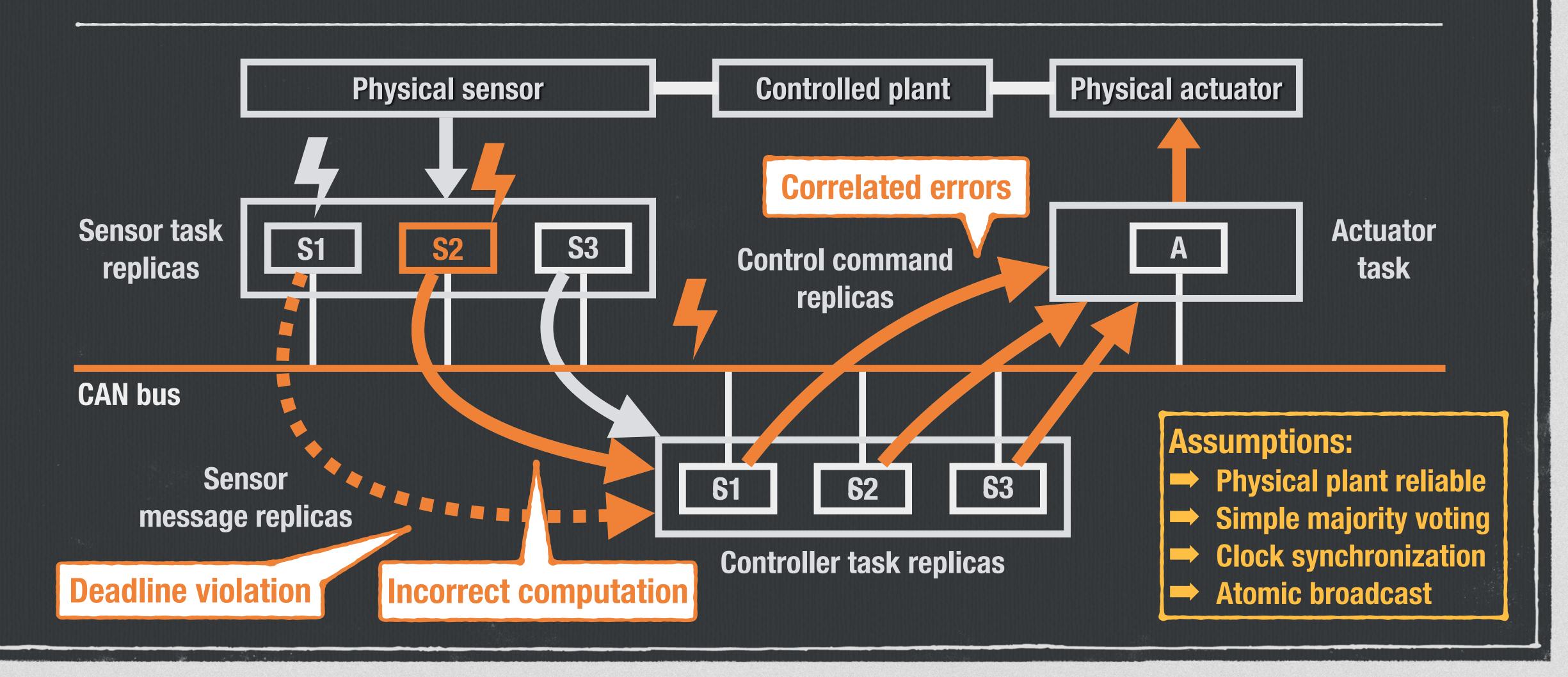
Fault tolerant single-input single-output (FT-SISO) networked control loop

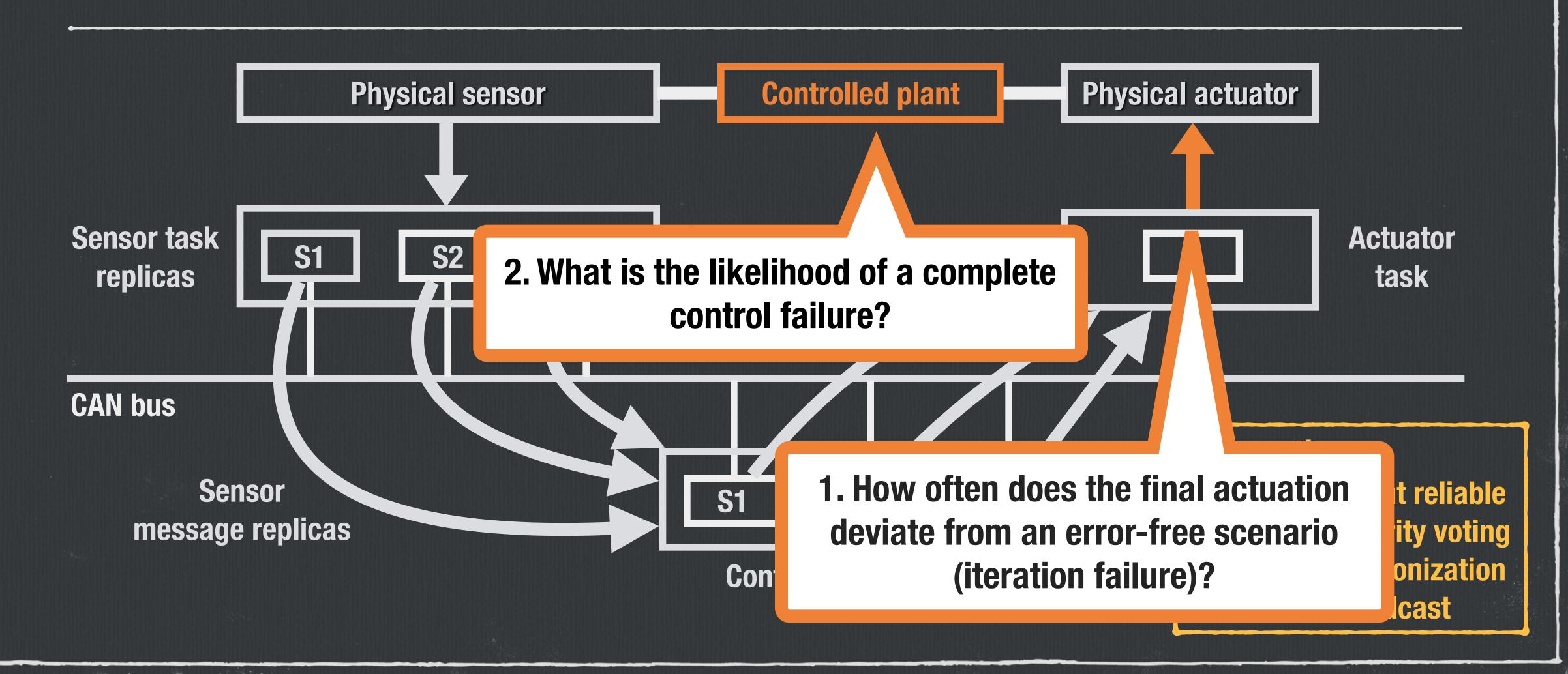








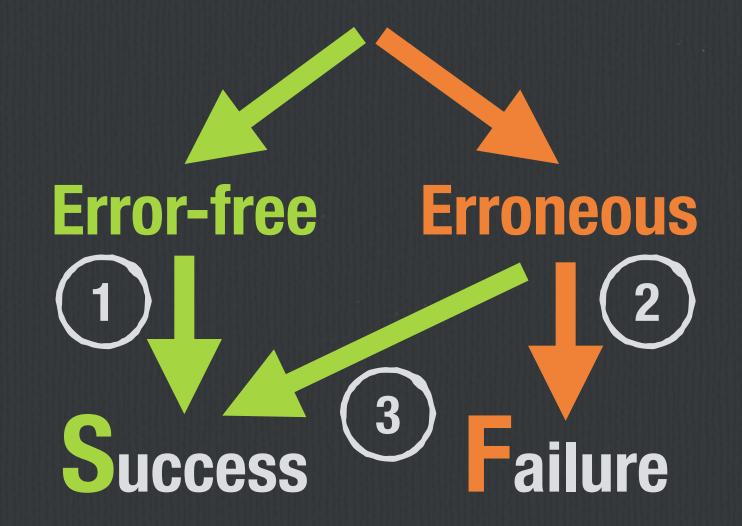




1. Modeling control loop iteration failures

Control loop iterations

- 1) Final actuation is successful
- 2) Final actuation failed (different from 1)
- Final actuation is successful (same as 1) despite the errors



Explicitly account for fault tolerance

2. Modeling control failure based on the (m, k)-firm constraint

Control loop iterations

Success Failure

time

SSSFSSSFSS

Hard constraint

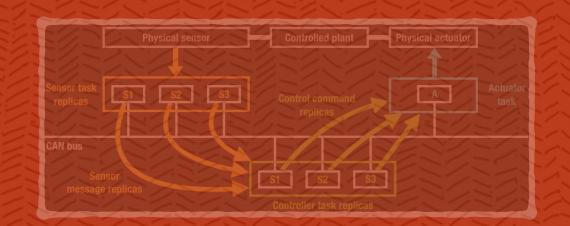
Control failure upon first iteration failure

(2, 3) constraint

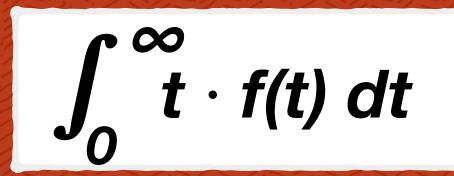
Control failure when less than 2 iterations successful in 3 consecutive iterations

Outline

Analysis of a Controller Area Network (CAN) based networked control system



System Model

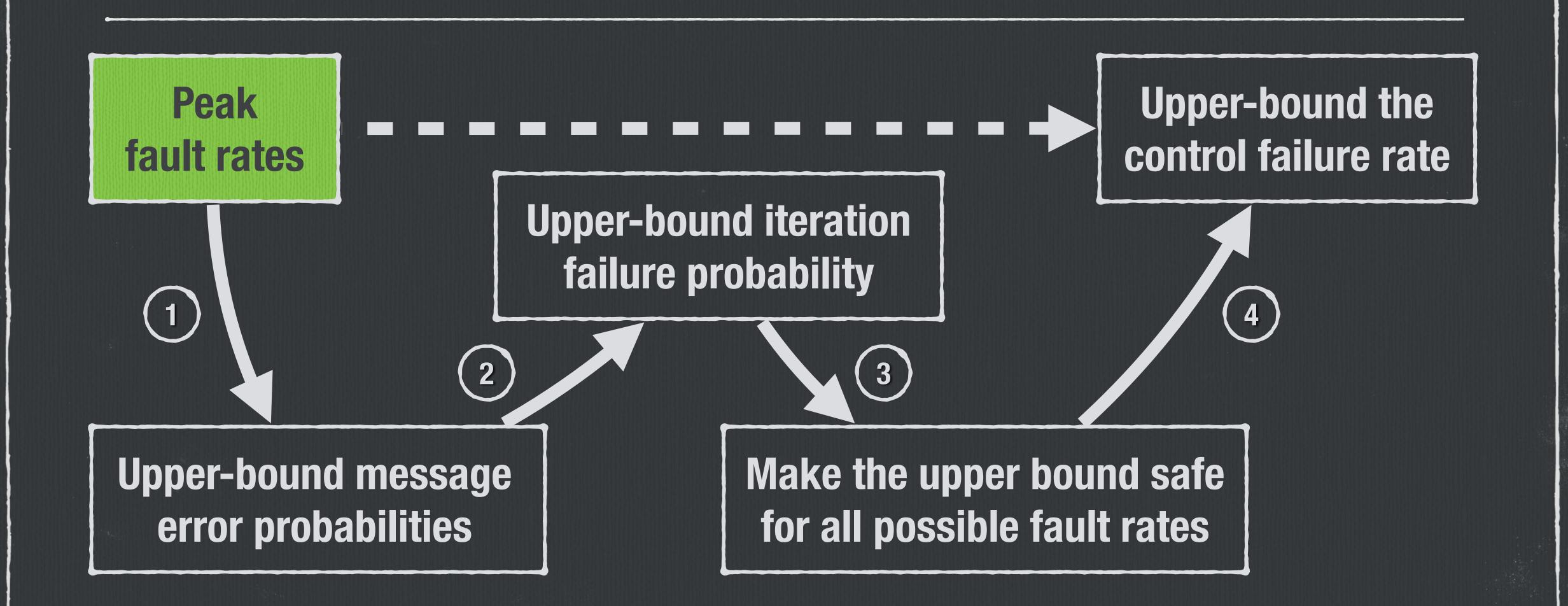


Analysis



Evaluation

Analysis steps



Upper-bounding the message error probabilities

Peak fault rates

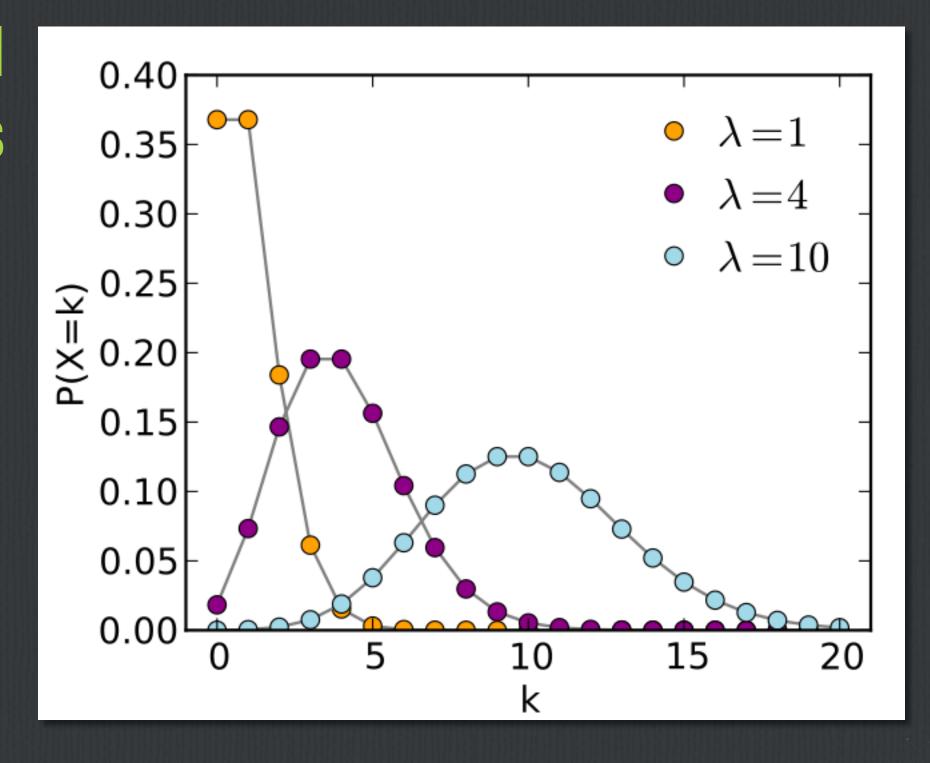
Using poisson model for fault arrivals

Based on the message parameters

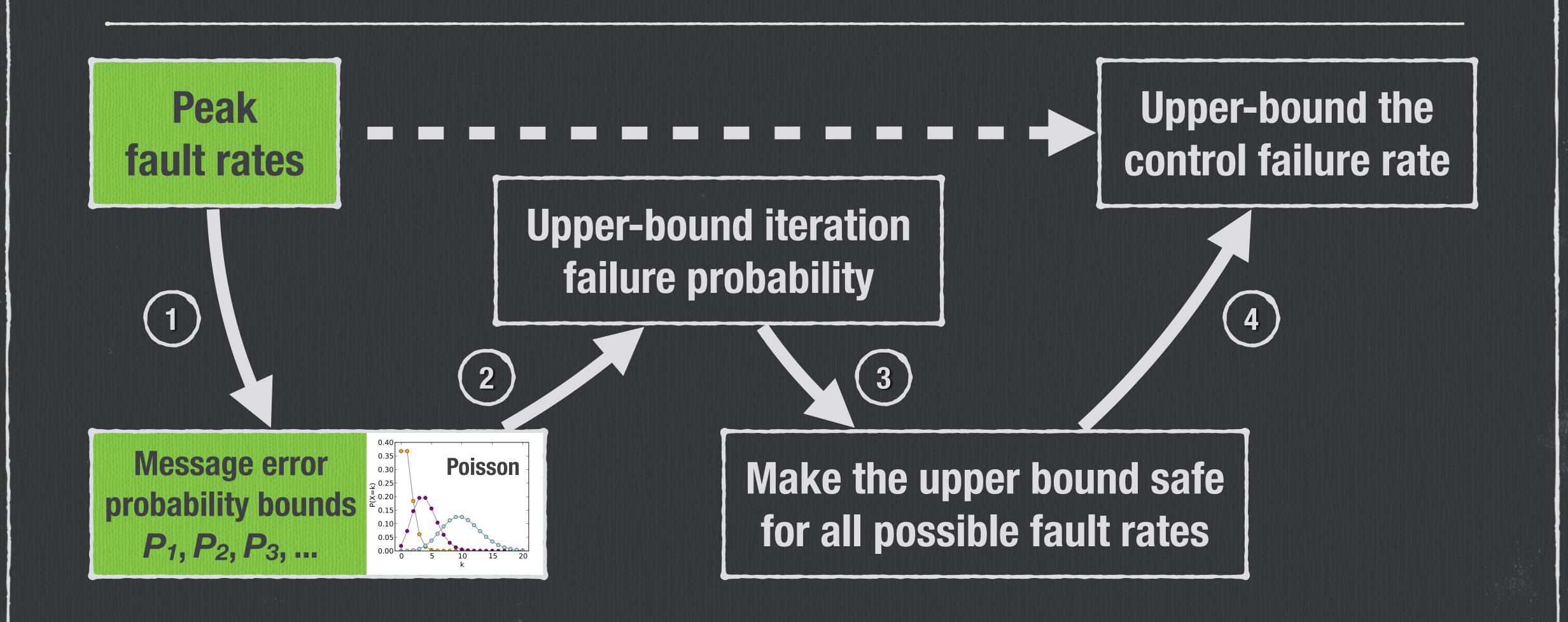
 $P_1 \ge P$ (msg. is omitted at time t)

 $P_2 \ge P$ (msg. is incorrectly computed)

 $P_3 \ge P$ (msg. is misses its deadline)



Analysis steps



Upper-bounding the iteration failure probabilities

Accounting for

- **→** all possible error scenarios
- error propagation and correlation
- voting protocol

Upper bounds on message error probabilities

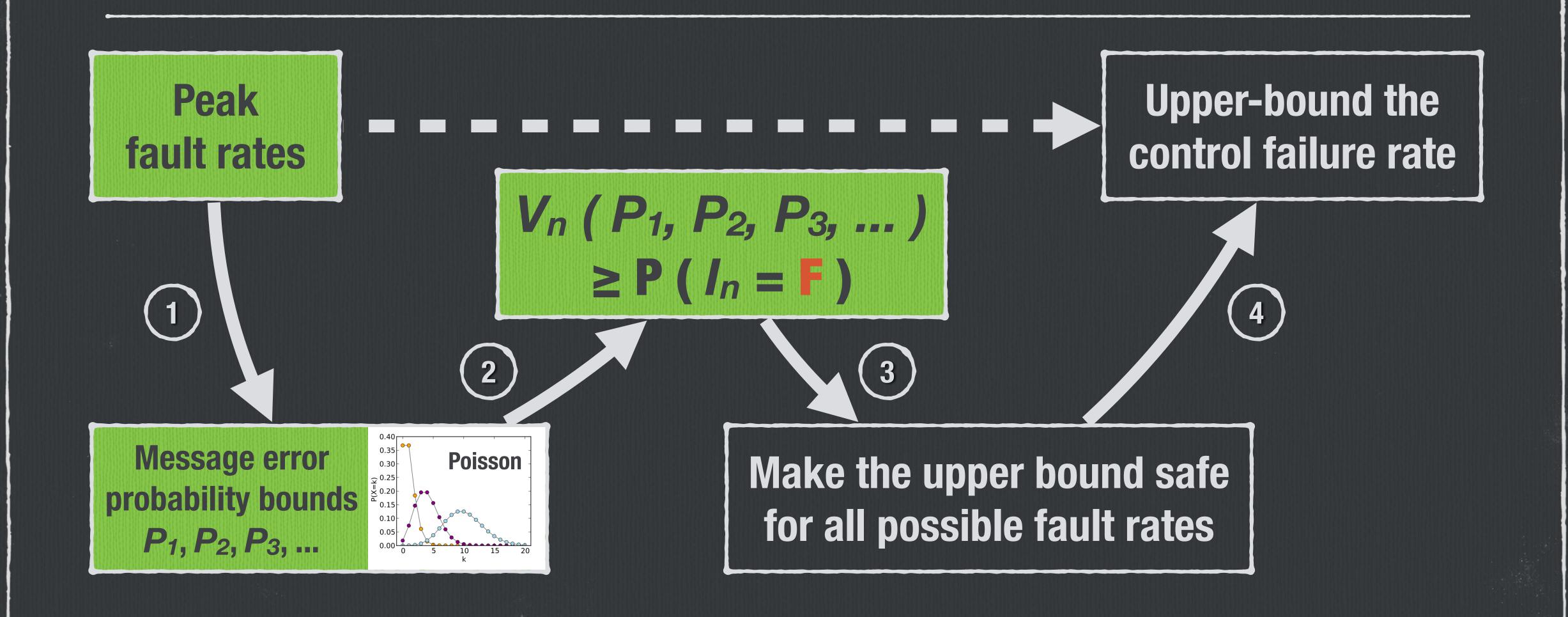
 $P_1 \ge P$ (msg. is omitted at time t)

 $P_2 \ge P$ (msg. is incorrectly computed)

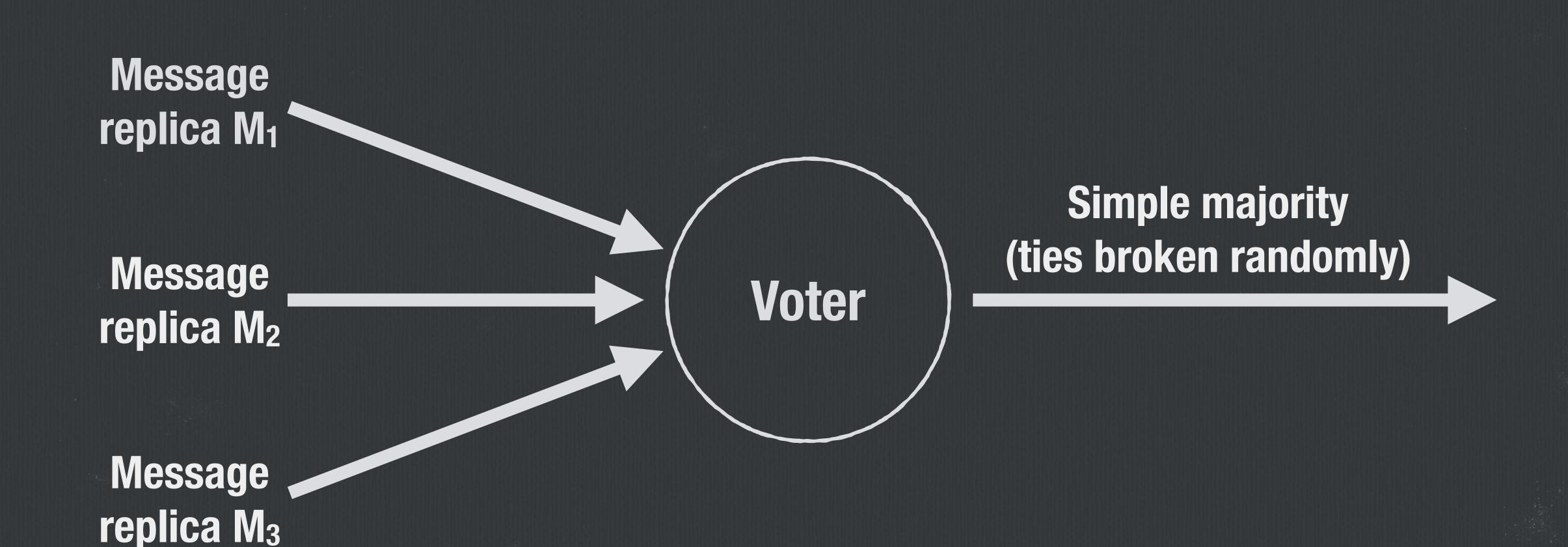
 $P_3 \ge P$ (msg. is misses its deadline)

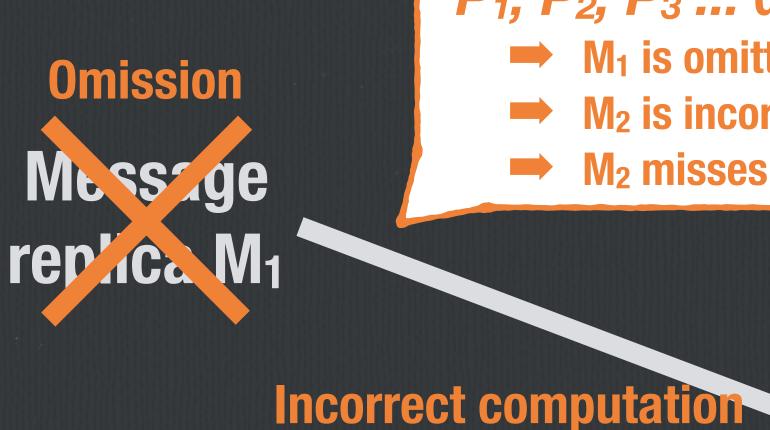
 $V_n(P_1, P_2, P_3, ...) \ge P(I_n = F)$

Analysis steps



Let's look at a simple example!





 P_1 , P_2 , P_3 ... defined such that:

- \rightarrow M₁ is omitted
- M₂ is incorrectly computed
- M₂ misses its deadline

Message & deadline violation Voter

Message replica M₃

replica M₂

Only M₃ participates in the voting process $V_n(P_1, P_2, P_3, ...) = 0$

Simple majority (ties broken randomly)



 P_1 , P_2 , P_3 ... defined such that:

- **→** M₁ is omitted
- **→** M₂ is incorrectly computed
- **→** M₂ misses its deadline

$$V_n(P_1, P_2, P_3, ...) = 0$$

Simple majority (ties broken randomly)

Incorrect computation

Message

A deadline violation

Voter

Message replica M₃

replica M₂

In practice, there may be no deadline violations!

→ The peak fault rates are just upper bounds

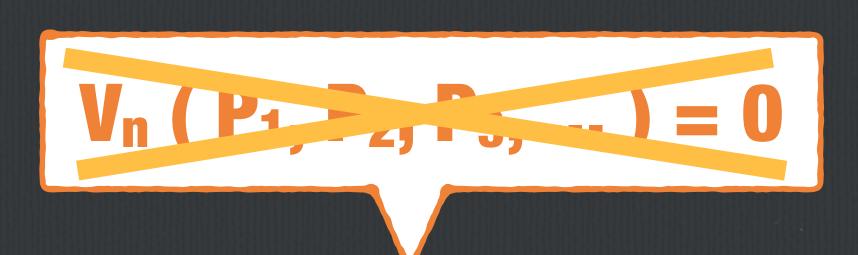


 P_1 , P_2 , P_3 ... defined such that:

→ M₁ is omitted

Incorrect computation

- **→** M₂ is incorrectly computed
- **→** M₂ misses its deadline



Simple majority (ties broken randomly)

Voter

 $V_n(P_1, P_2, P_3, ...) = 0.5$

Message replica M₃

Message

replica M₂

In practice, there may be no deadline violations!

→ The peak fault rates are just upper bounds

$$V_n(P_1, P_2, P_3, ...) \ge P(I_n = F)$$

+

A fudge factor ∆ is added to ensure monotonicity*

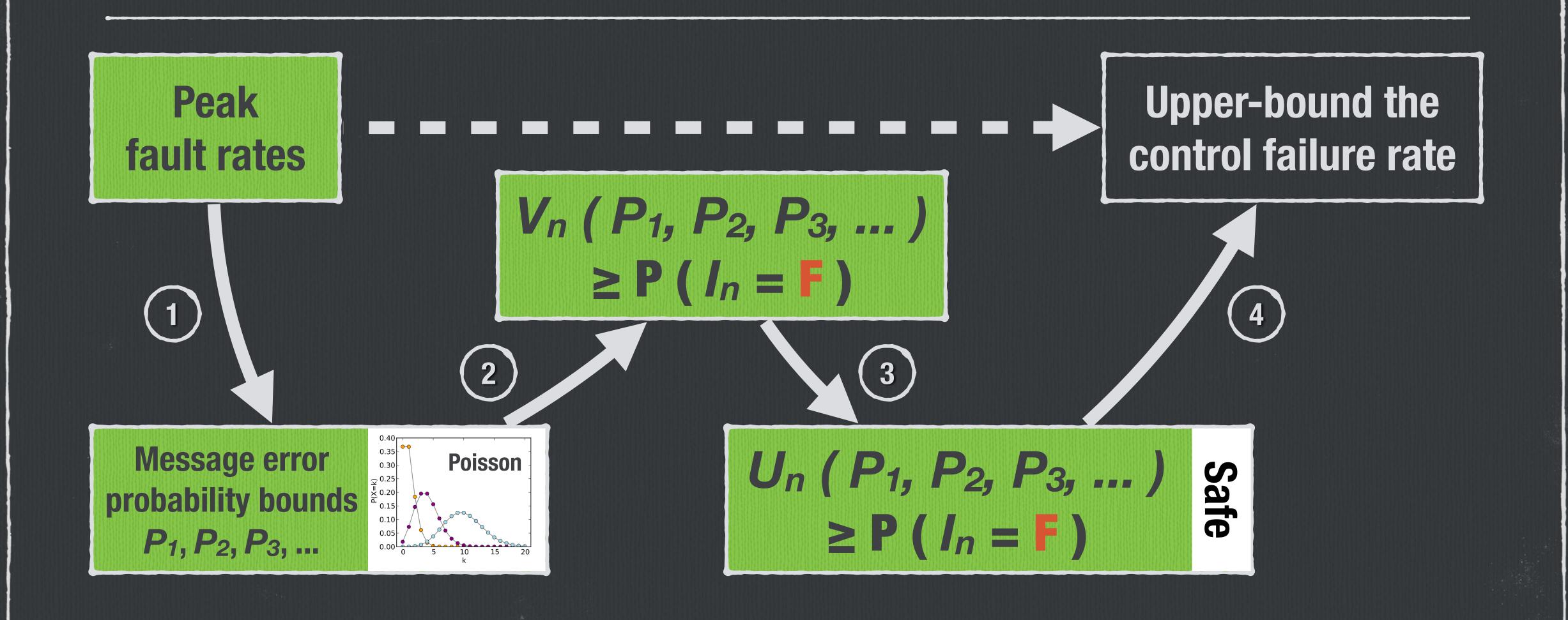
П

$$|U_n(P_1, P_2, P_3, ...) \ge P(I_n = F)|$$

Safe if V_n is monotonic in P_1 , P_2 , P_3 , ...

*Arpan Gujarati, Mitra Nasri, and Björn B Brandenburg. Quantifying the resiliency of fail-operational real-time networked control systems. Technical Report MPI-SWS2018-005, Max Planck Institute for Software Systems, Germany, 2018. URL: http://www.mpi-sws.org/tr/2018-005.pdf.

Analysis steps



Upper-bounding the control failure rate (Failures-In-Time or FIT)

$$U_n(P_1, P_2, P_3, ...)$$

 $\geq P(I_n = F)$

Using prior work*

Scalable and numerical, but sound, analysis

FIT
(expected # failures in 1 billion hours)

= 10⁹ / MTTTF (in hours)
(Mean Time To first control Failure)

(probability density function)

(probability density function)

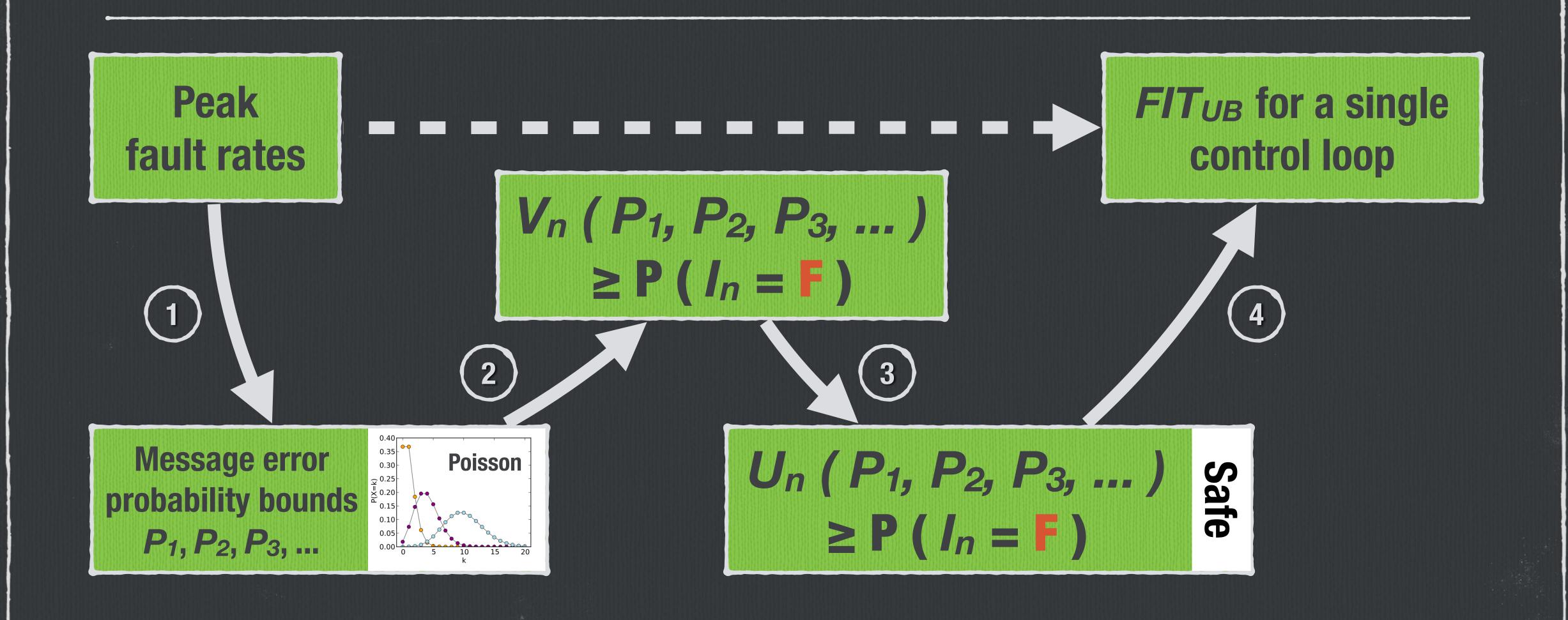
f(t) = P (first control failure at time t)

= P (first violation of (2, 3)-firm constraint at time t)

= P (first instance of FSF | FFS | SFF | FF at time t)

*M. Sfakianakis, S. Kounias, and A. Hillaris. "Reliability of a consecutive k-out-of-r-from-n: F system." IEEE Transactions on Reliability 41, no. 3 (1992): 442-447.

Analysis steps

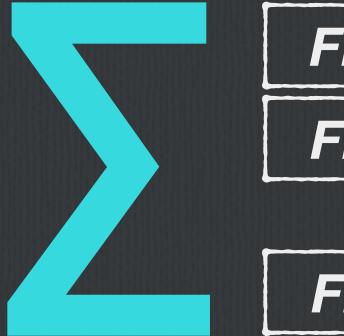


Analysis steps

Peak fault rates

FITUB for a single control loop

Upper bound on the FIT rate of the entire networked control system



FITUB for L₁

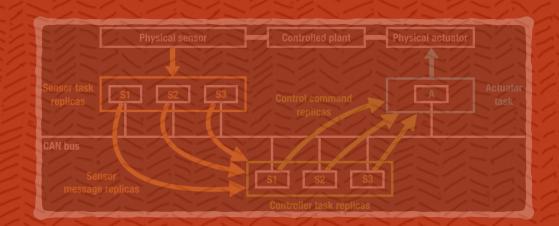
FIT_{UB} for L₂

FIT_{UB} for L_n

Compute FIT bounds for all control loops in the networked control system

Outline

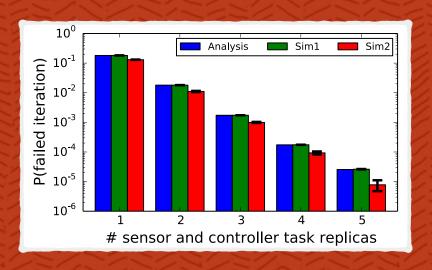
Analysis of a Controller Area Network (CAN) based networked control system



System Model



Analysis



Evaluation

Evaluation overview

- ☐ How accurate is the analysis?
 - Comparison with simulation results
- ☐ Case study: FIT vs. (m, k) constraints vs. replication schemes

CAN-based active suspension workload*

- ☐ Four control loops L₁, L₂, L₃, L₄
 - to control the four wheels with magnetic suspension

This talk: Control loop L_1 's tasks were replicated

In the paper: Experiments with <u>all</u> replica schemes

Messages	Length	Period (ms)	Deadline (ms)	Priority
Clock sync.	1	50	50	High
Current mon.	1	4	4	
Temperature	1	10	10	
L ₁ messages	3	1,75	1,75	
L ₂ messages	3	1,75	1,75	
L ₃ messages	3	1,75	1,75	
L ₄ messages	3	1,75	1,75	
Logging	8	100	100	Low

^{*}Adolfo Anta and Paulo Tabuada. On the benefits of relaxing the periodicity assumption for networked control systems over CAN. In Proceedings of the 30th Real-Time Systems Symposium, pages 3–12. IEEE, 2009.

How accurate is the analysis?

Iteration failure probability bound

$$U_n(P_1, P_2, P_3, ...) \ge P(I_n = F)$$

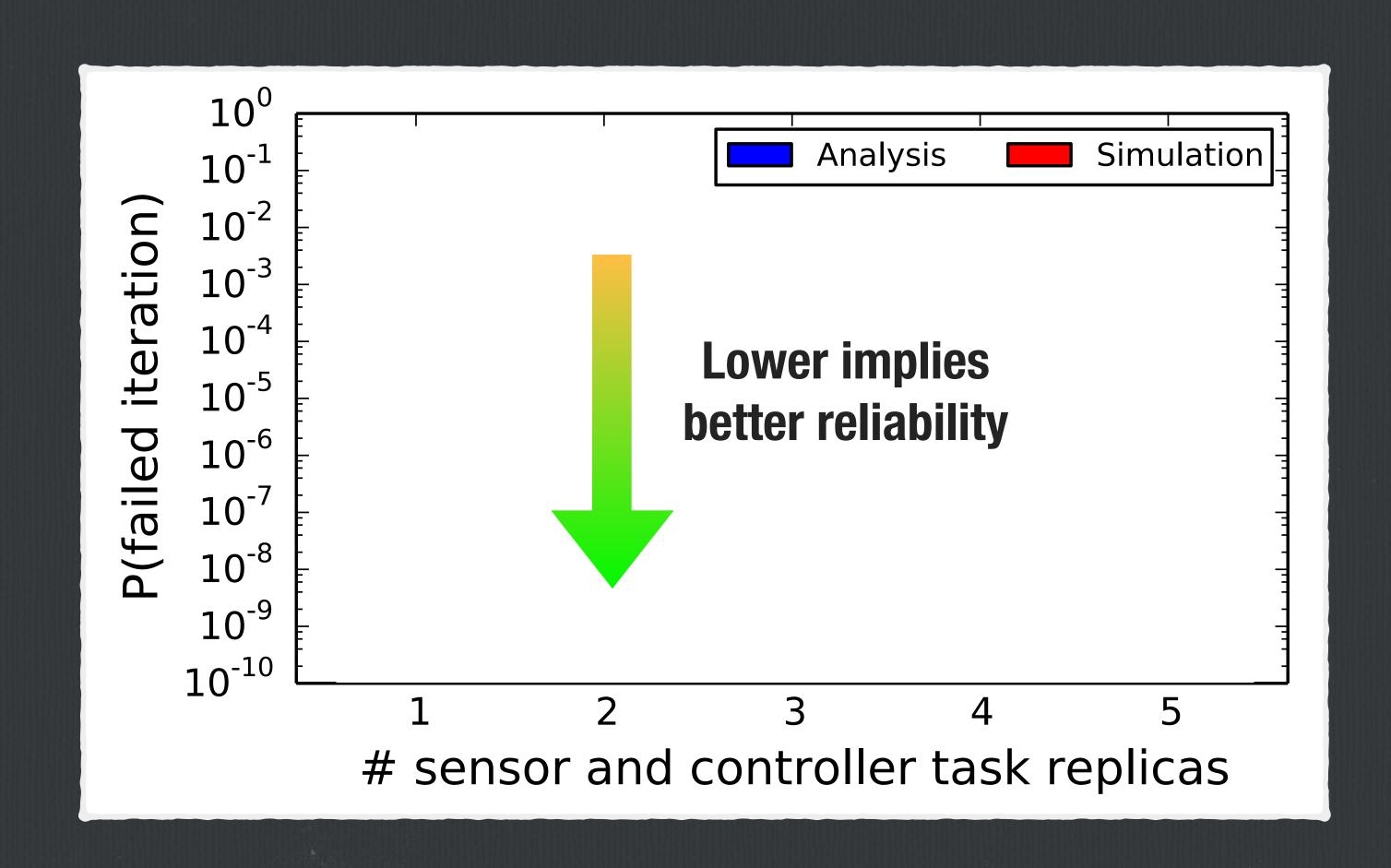
Simulation is not safe

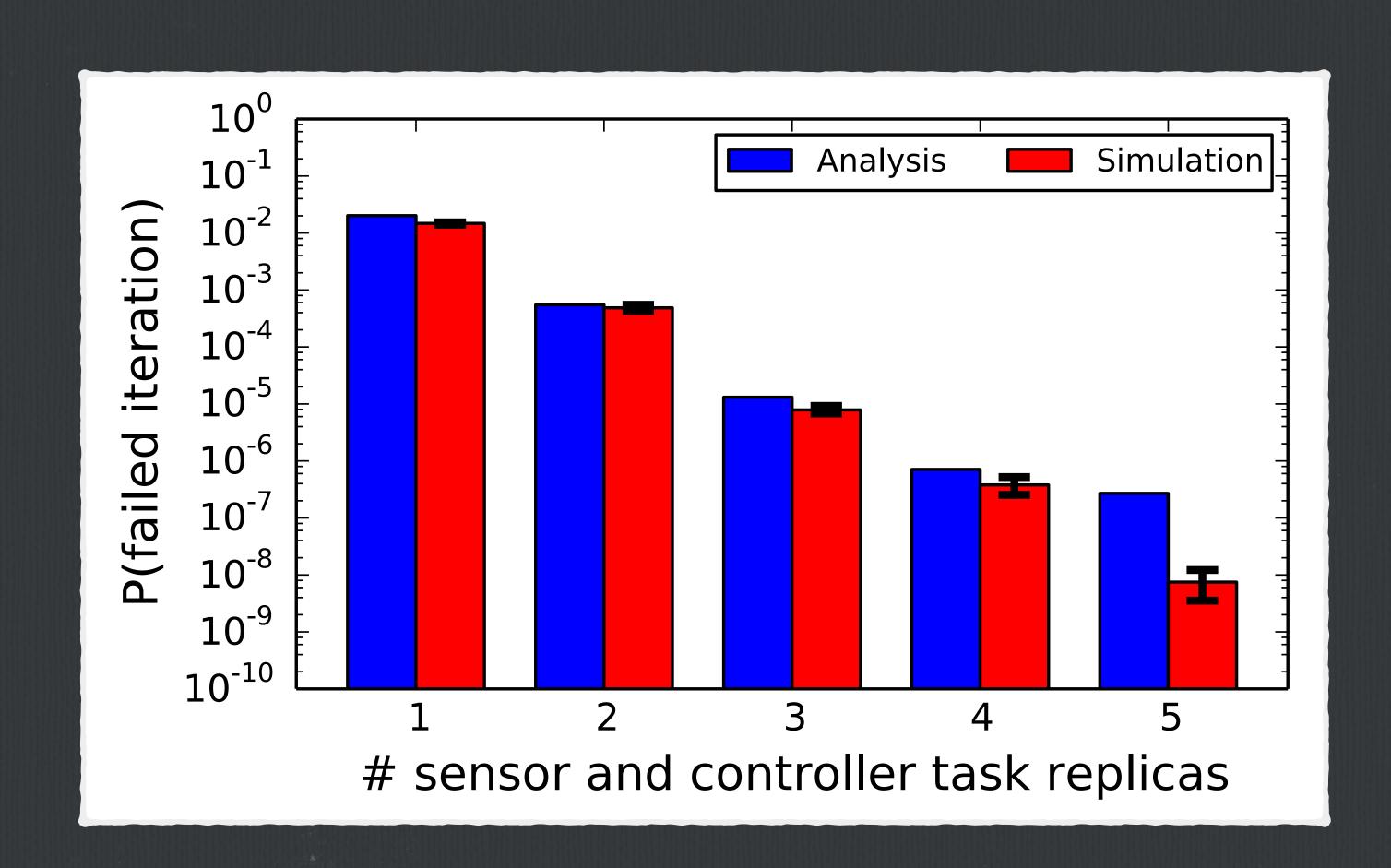
Discrete event simulation of a CAN-based system

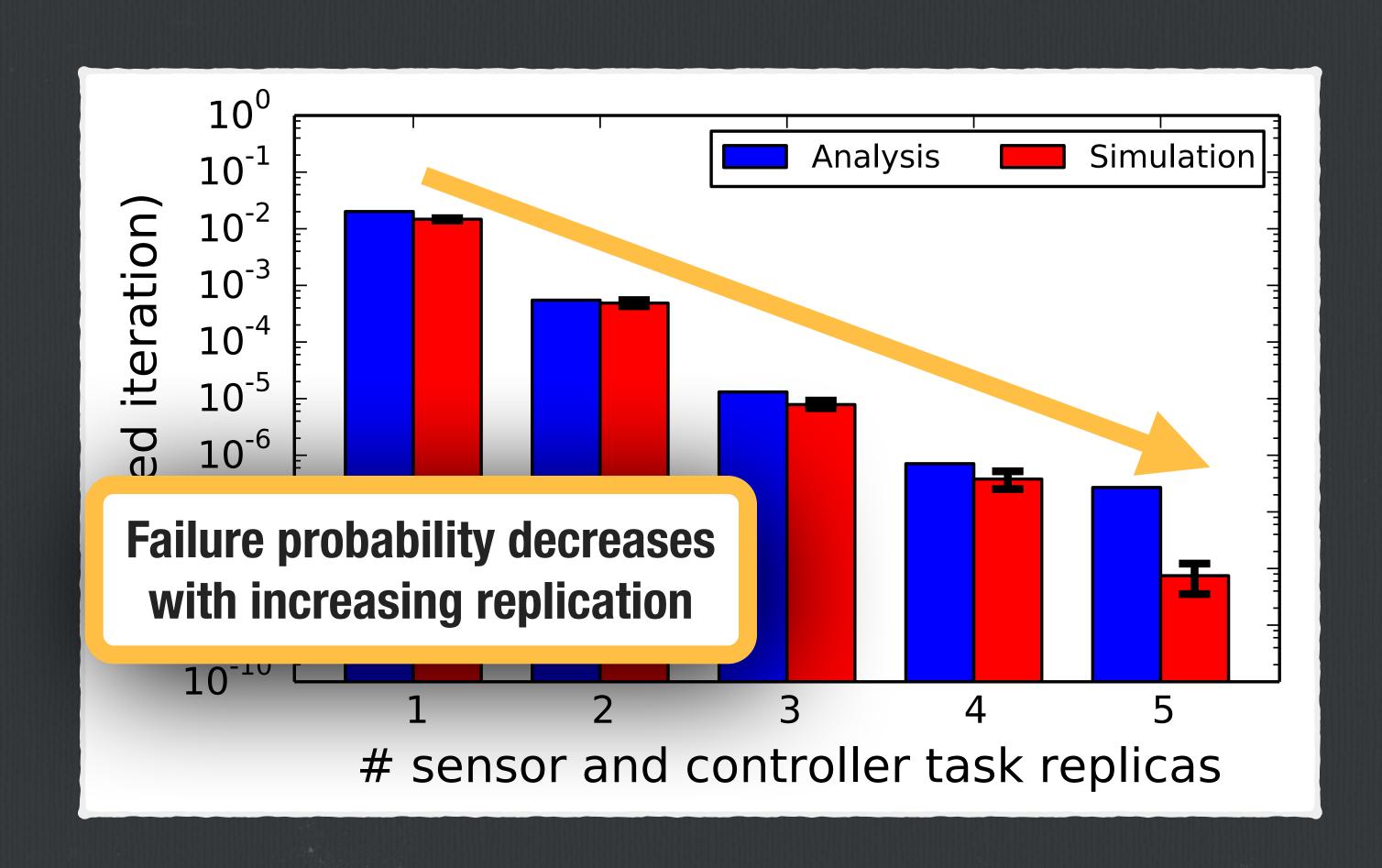
Poisson process for CAN bus faults

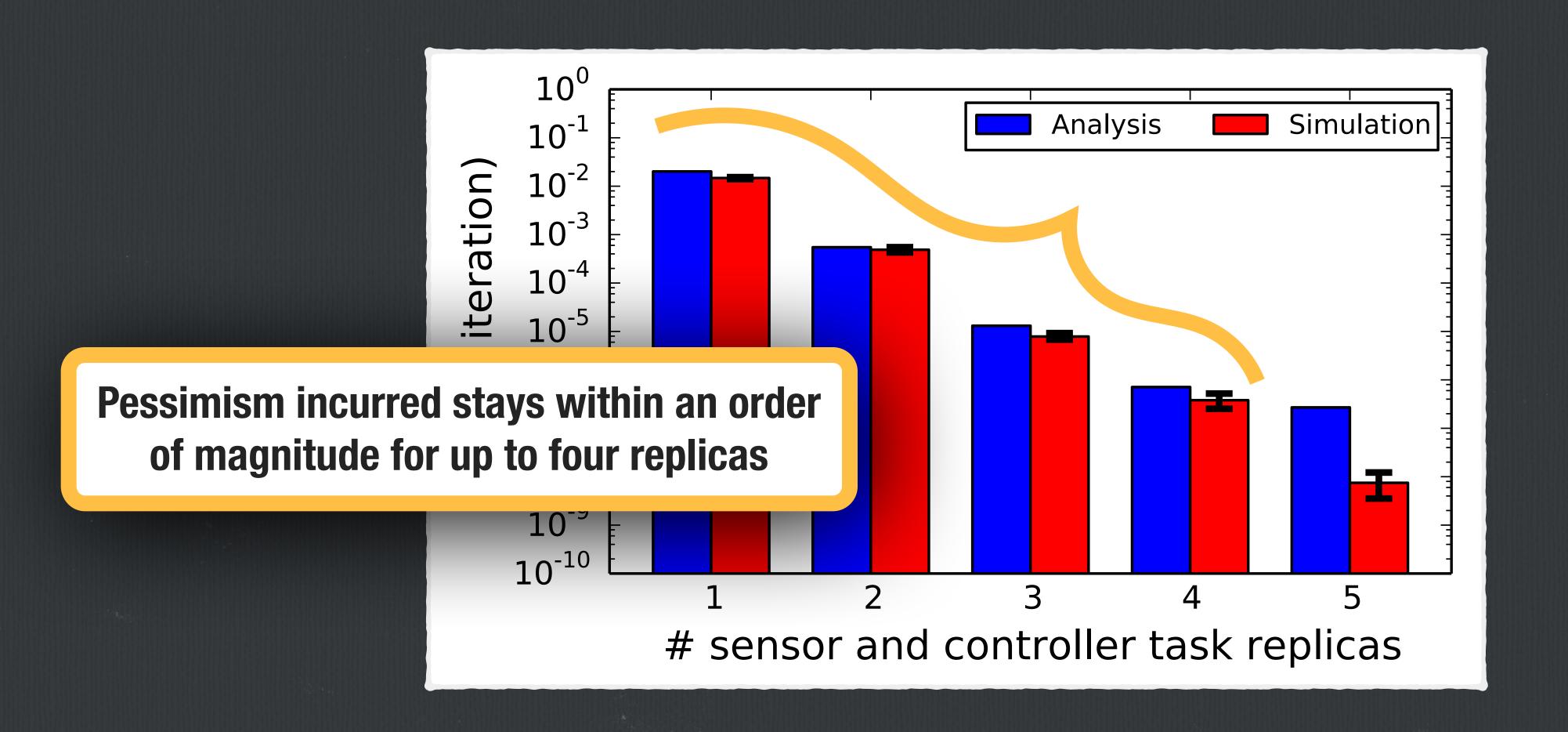
Poisson process for faults on Host 1

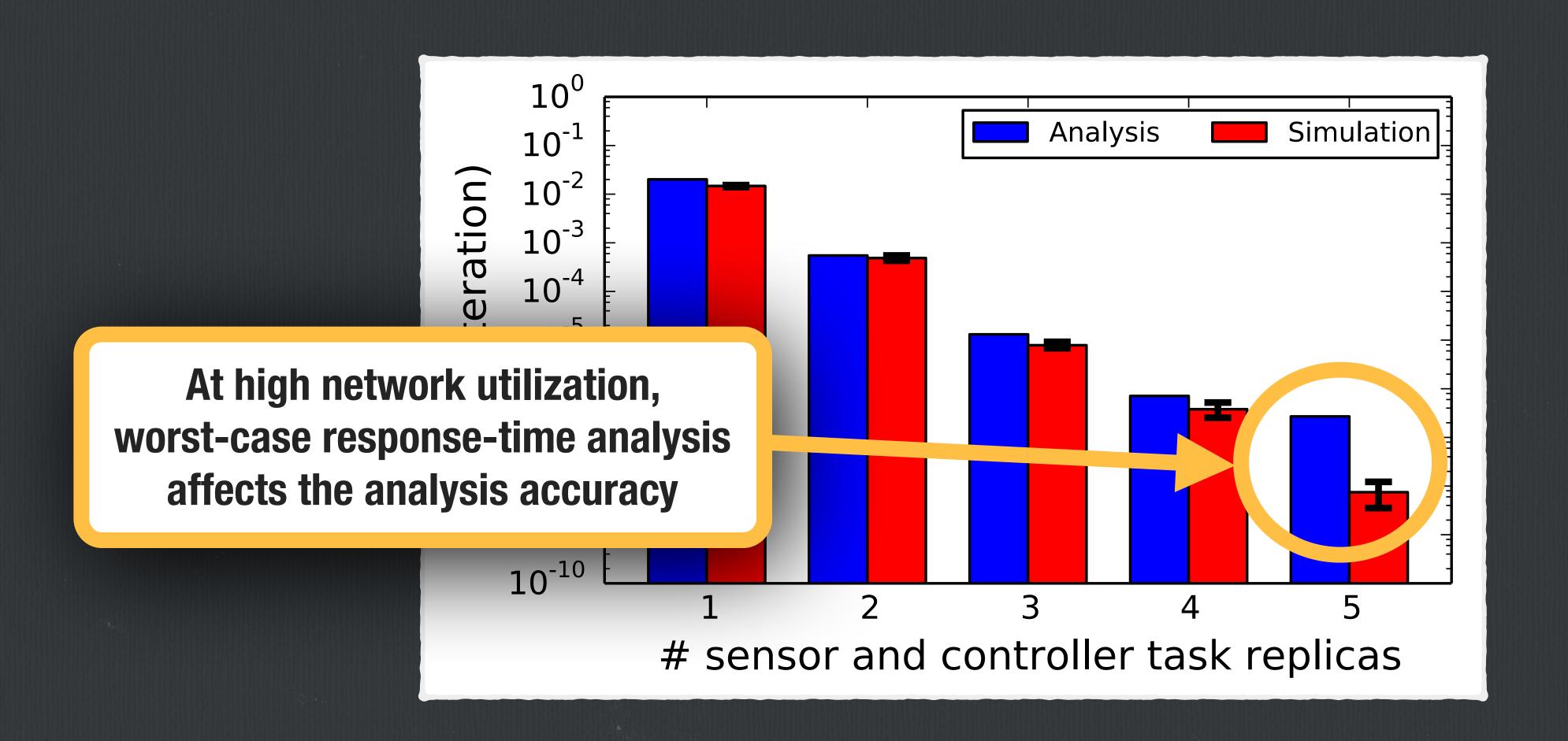
... and so on





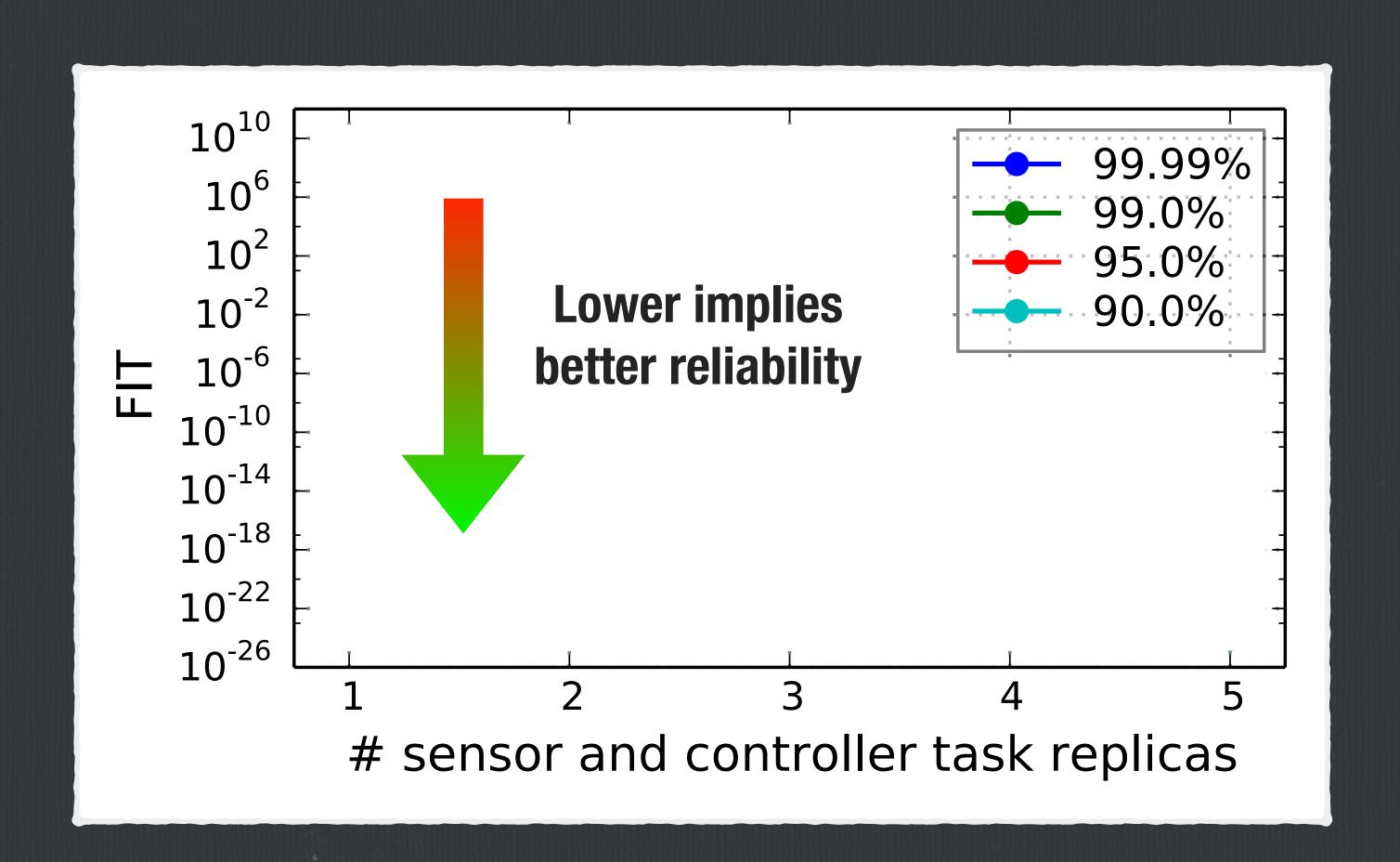


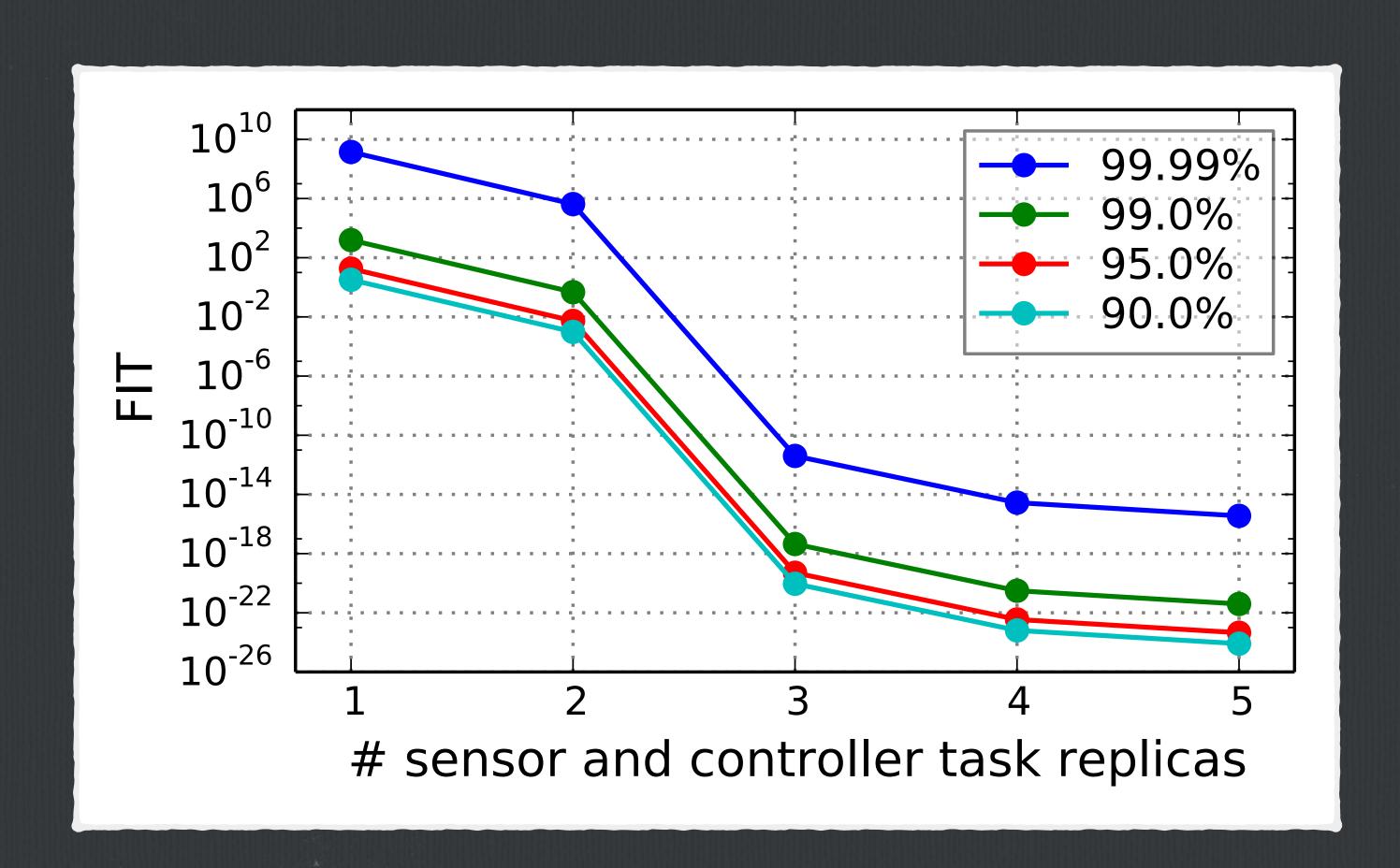


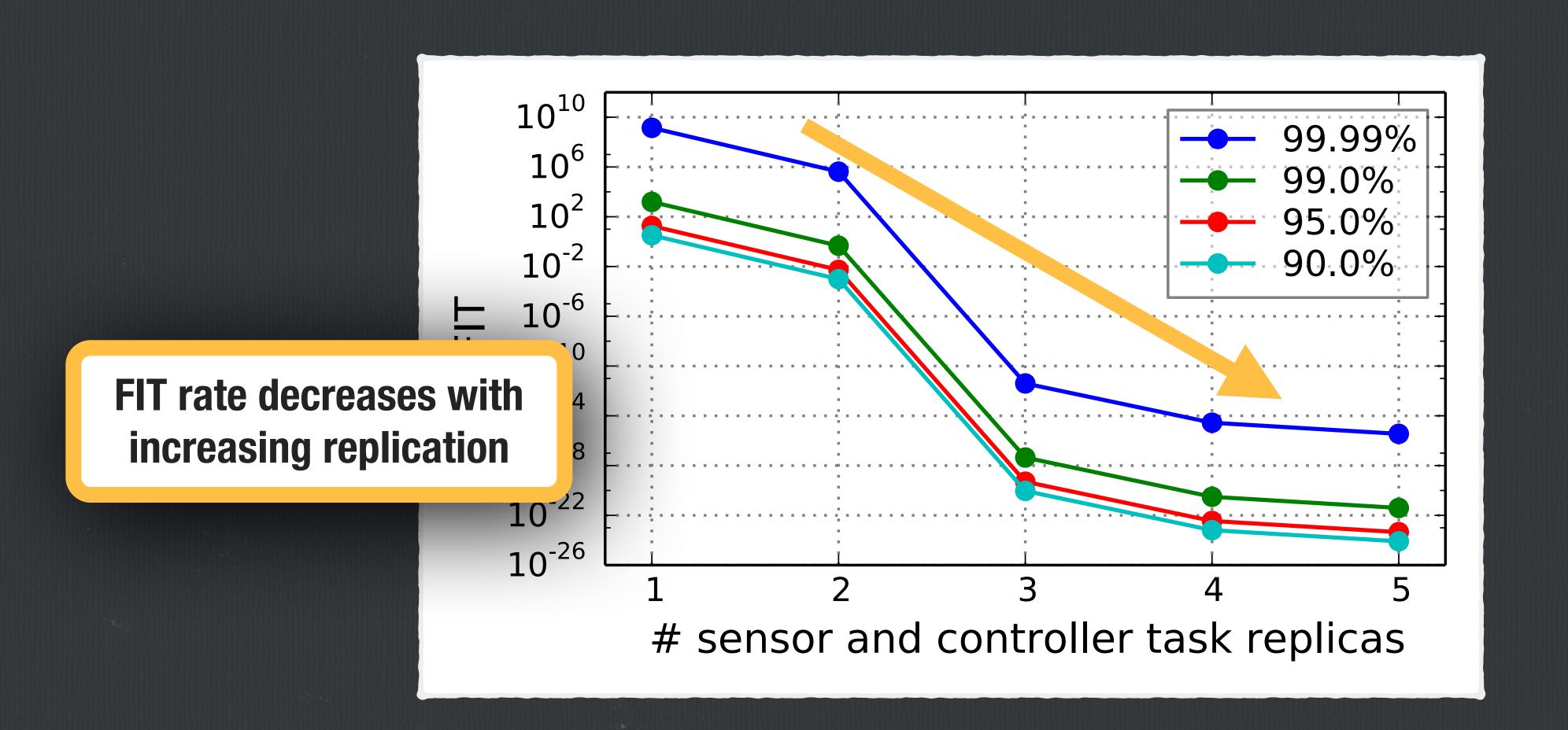


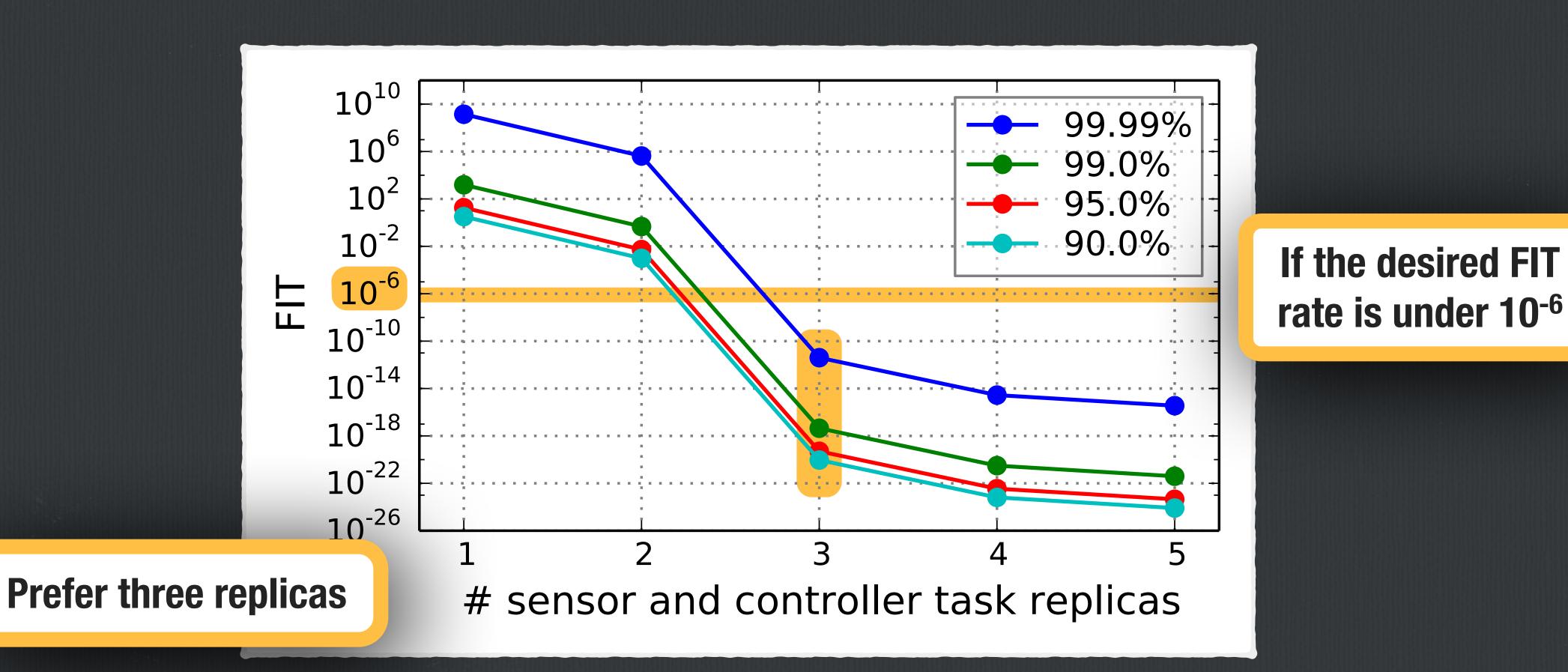
Case study

- ☐ FIT analysis for different (m, k)-firm constraints
 - **→** (9, 100) ~ 9%
 - **→** (19, 20) ~ 95%
 - **→** (99, 100) ~ 99%
 - **→** (9999, 10000) ~ 99.99%
- ☐ Replication factor of loop L₁'s tasks varied from 1 to 5
- ☐ What should be the replication factor to achieve FIT under 10⁻⁶?









Summary

- ☐ A safe Failures-In-Time (FIT) analysis for networked control systems
 - **→ CAN-based networked control system model**
- ☐ Focus on failures and errors due to transient faults
 - omission errors
 - incorrect computation errors
 - **→** transmission errors

Future work: Byzantine errors + BFT protocols

- ☐ ... and on robust systems that can tolerate a few iteration failures
 - (m,k)-firm model for control failure

Accounting for other robustness criteria