

INSTITUTE
OF COMMUNICATION,
INFORMATION
AND PERCEPTION
TECHNOLOGIES



Scuola Superiore
Sant'Anna



Need for Reservation Servers with Constrained Deadlines

Daniel Casini, Alessandro Biondi, and Giorgio Buttazzo

Scuola Superiore Sant'Anna – ReTiS Laboratory

Pisa, Italy

Why using constrained-deadlines?

Recent work showed that **Semi-Partitioned scheduling** can achieve **high schedulability** performance:

- “*Global Scheduling Not Required*” by Brandenburg and Gul for **static workloads** (RTSS 2016)
- “*Semi-Partitioned Scheduling of Dynamic Real-Time Workload*” by Casini et al. for **dynamic workloads**
(29th June, 15:30 PM @ ECRTS 2017)

Why using constrained-deadlines?

- Supporting constrained-deadlines is an open problem also for the **SCHED_DEADLINE** scheduling class of **Linux** (based on reservation with the **CBS** algorithm)
- Currently discussed also in the **Linux kernel mailing list**

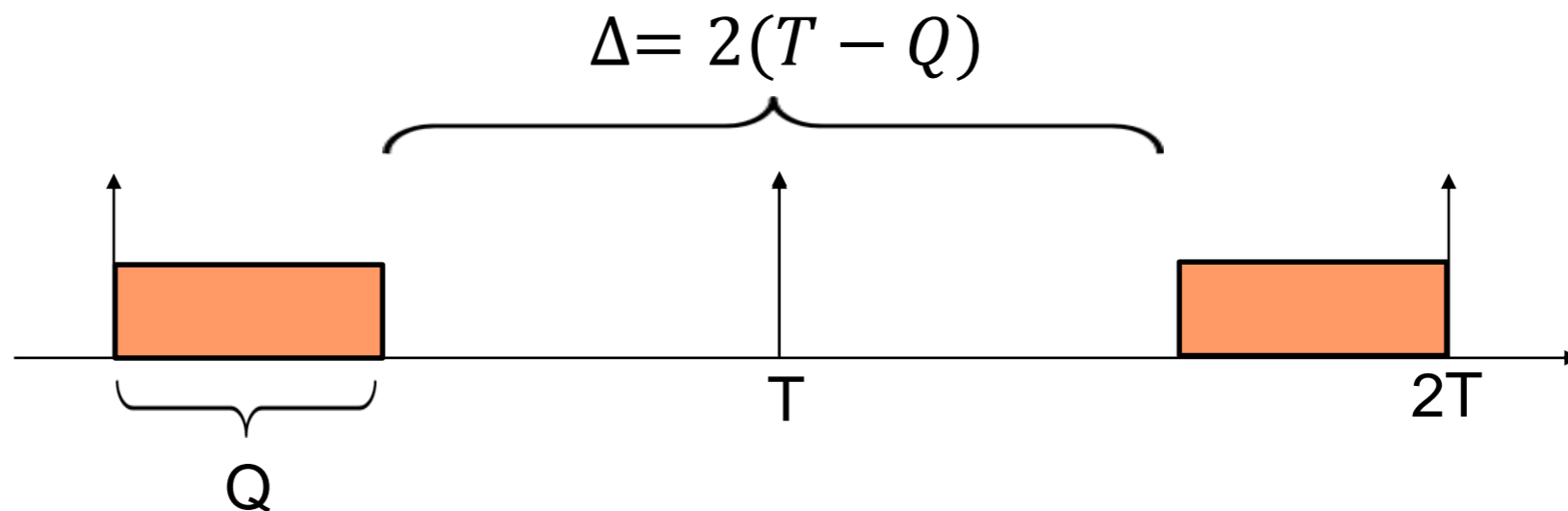


Linux

Hard Constant Bandwidth Server

- H-CBS is a reservation algorithm allowing to guarantee:
- A bandwidth $\alpha = \frac{Q}{T}$
- A bounded **maximum** service-delay $\Delta = 2(T - Q)$

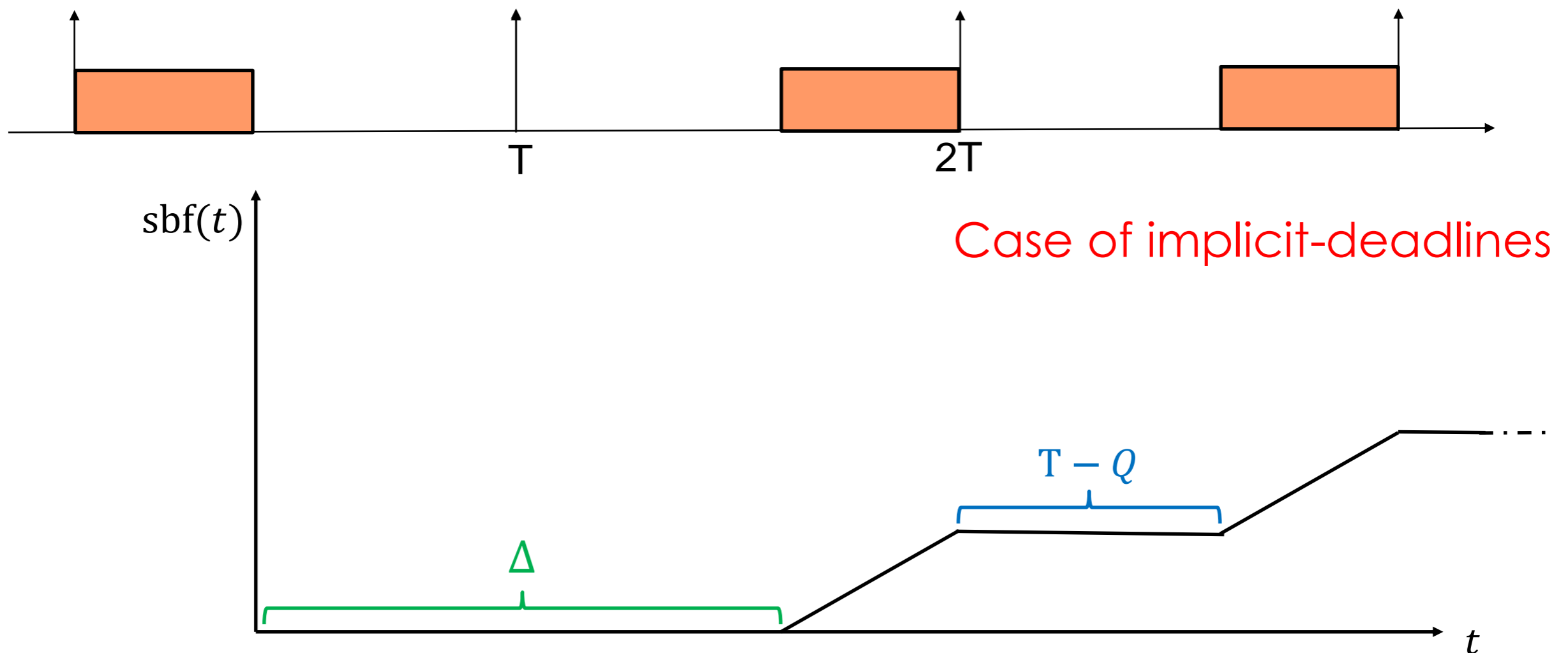
Worst-case scenario
for the service delay



Used in several works and implementations

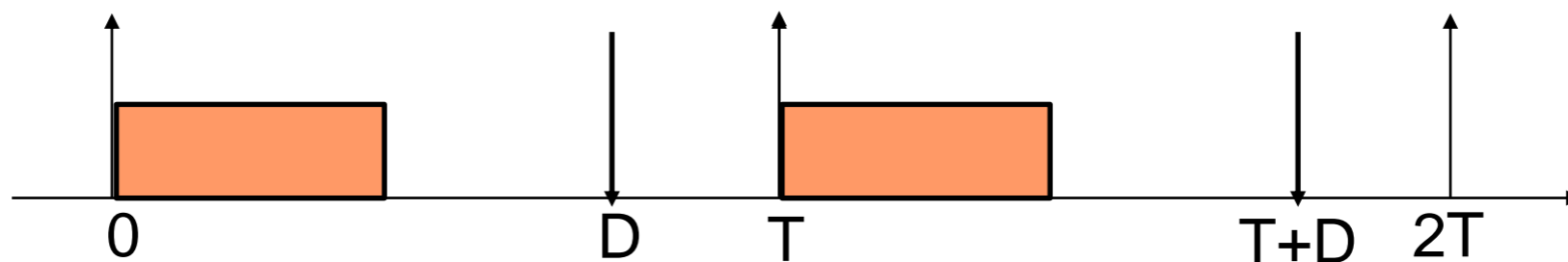
Importance of a bounded delay

A **bounded-delay** allows deriving a **supply function** that can be used for testing the **schedulability** of the workload running inside the server:



H-CBS and constrained-deadlines

- As long as the server behaves (in the worst-case) as a **standard periodic/sporadic task** with constrained deadlines, **existing EDF schedulability theory** can be applied
- The core issue is how to guarantee that the **demand** generated by the server **never exceeds** the one of a corresponding sporadic task in **all possible scenarios...**



H-CBS key rule

□ H-CBS has a specific rule when the server wakes up from the idle state:

➤ **Rule 2:** “When H-CBS is idle and a job arrives at time t , a *replenishment time* is computed as $t_r = d - \frac{q}{\alpha}$ ”

- Then, if $t < t_r$ the server is **suspended** until time t , where the budget is **replenished** and the absolute **deadline** is **postponed** to time $t_r + T$;
- otherwise, the budget is **immediately replenished** and the absolute deadline is postponed to $t + T$.

H-CBS and constrained-deadlines

□ H-CBS rules are not directly applicable in case of constrained-deadlines:

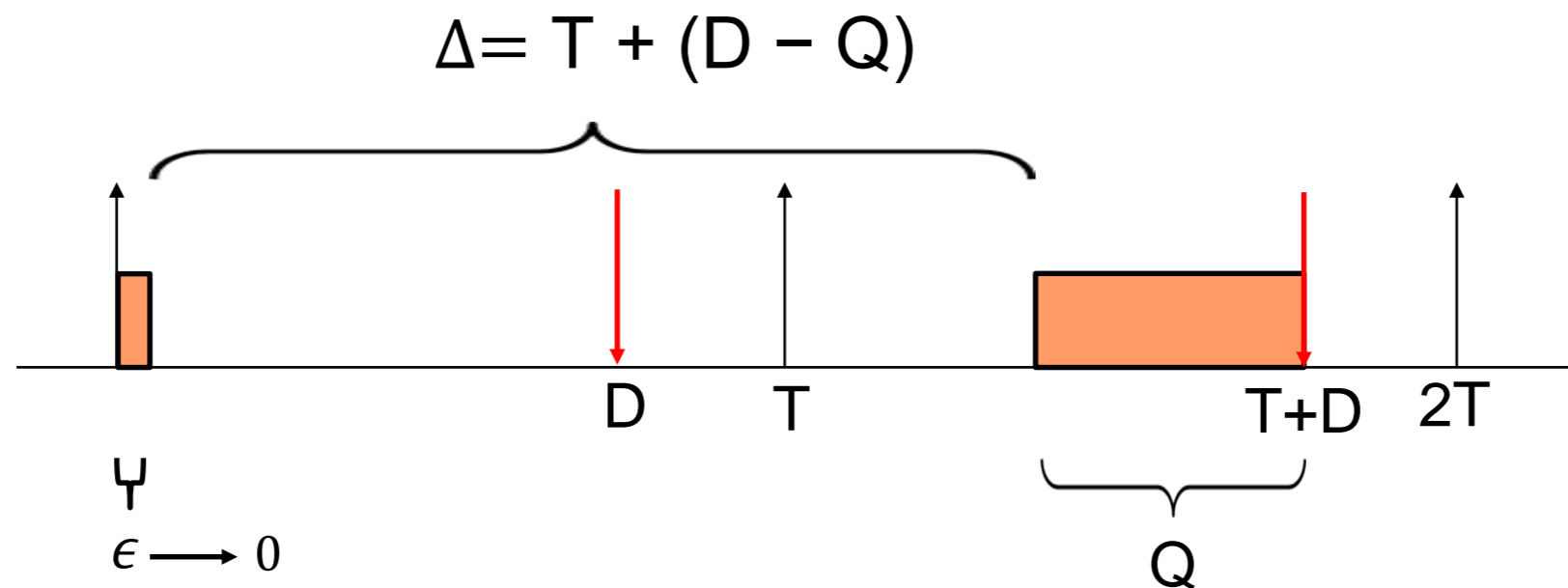
➤ **Rule 2:** “When H-CBS is idle and a job arrives at time t , a **replenishment time** is computed as $t_r = d - \frac{q}{\alpha}$ ”

This rule has been derived by **EDF schedulability** theory for **implicit-deadline tasks** (utilization-based), which indeed cannot be re-used to ensure **schedulability** with **constrained deadlines**!

Naïve solution

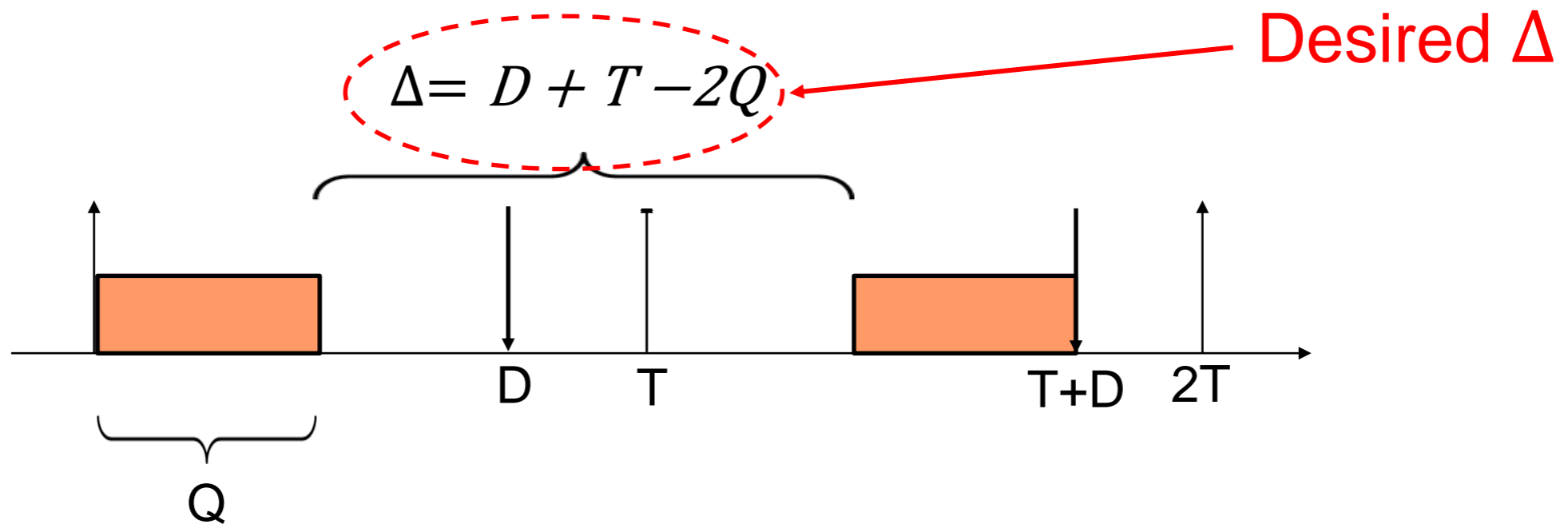
Mimic the polling server

- **New Rule:** “When H-CBS goes IDLE, *discard* all the budget. The budget is *replenishment* only at server periods, i.e., $t_r = kT_i$ ”



The worst-case service **delay** is **much higher!**

Questions



How to modify the **replenishment rules** for obtaining a **better** maximum-service **delay**?

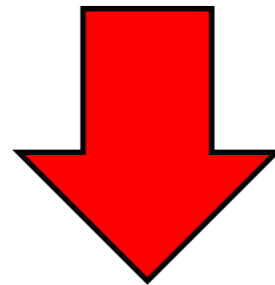


Is it possible to achieve a maximum service **delay** equal to $\Delta = D + T - 2Q$?

Issues with shared resources

BROE

- ✓ Avoids budget overruns
- ✓ Ensures bandwidth isolation
- ✓ Guarantees bounded-delay

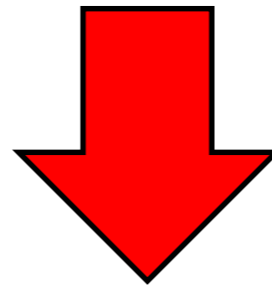


The protocol is based on a **proportional deadline-postponement** rule which relies on the **server bandwidth** (again, EDF schedulability theory for implicit-deadlines)

Issues with shared resources

BROE

- ✓ Avoids budget overruns
- ✓ Ensures bandwidth isolation
- ✓ Guarantees bounded-delay



How to guarantee a **bounded-delay** partition
in the presence of **shared-resources**?

Issues with admission control

- Replenishment rules are based on the admission test, so another question arise:



Which admission control test should be used for admitting reservations?

- We expect that the adopted admission test will strongly influence the server rules
- An efficient (and hence possibly sufficient) admission test would also reduce the server run-time overhead

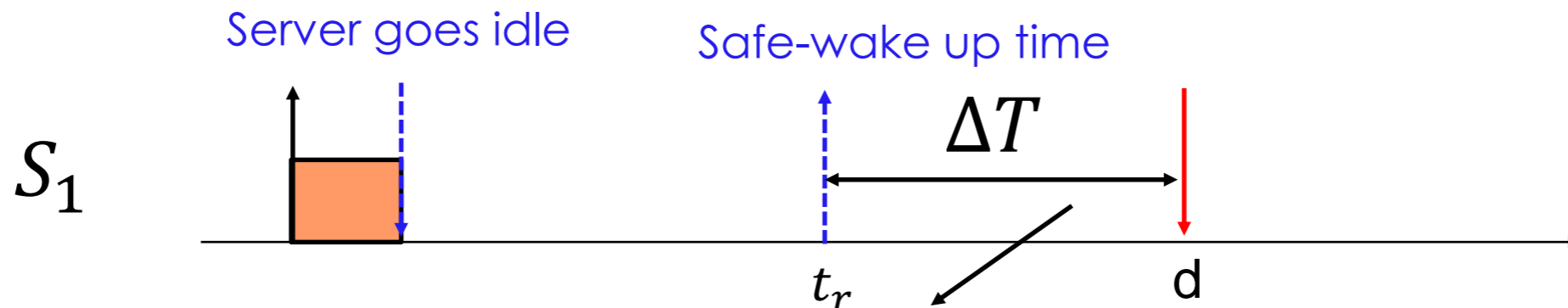
Issues with admission control

- With implicit-deadline the admission test of the H-CBS (based on EDF) is very simple:

$$\sum \alpha_i \leq 1$$

Exact test
Constant-time complexity

- This is relevant to our purpose because the H-CBS rule builds upon the schedulability test



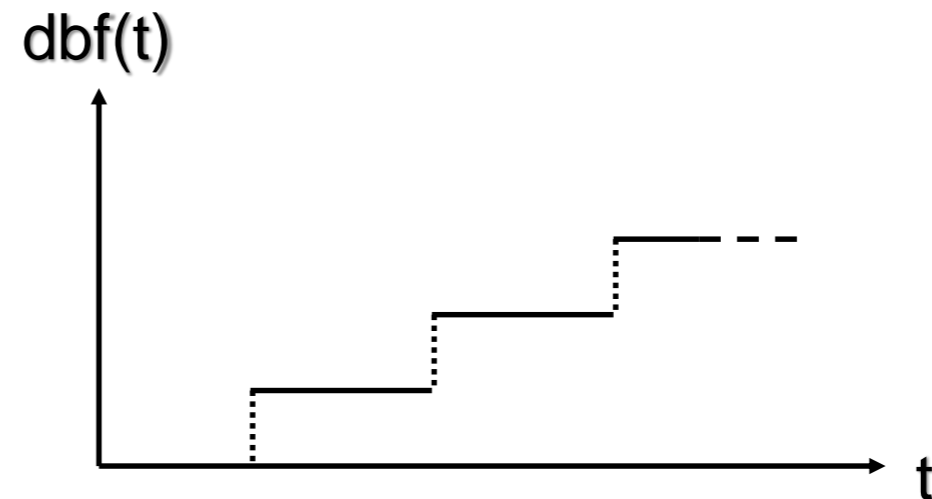
What is the t_r which guarantees a bandwidth α in ΔT ?

$$\alpha \Delta T = q \rightarrow \alpha(d - t_r) = q \rightarrow t_r = d - \frac{q}{\alpha}$$

Issues with admission control

- Conversely, considering constrained-deadlines the **schedulability check** is based on **Processor Demand Criterion** (Baruah et al. 1990)

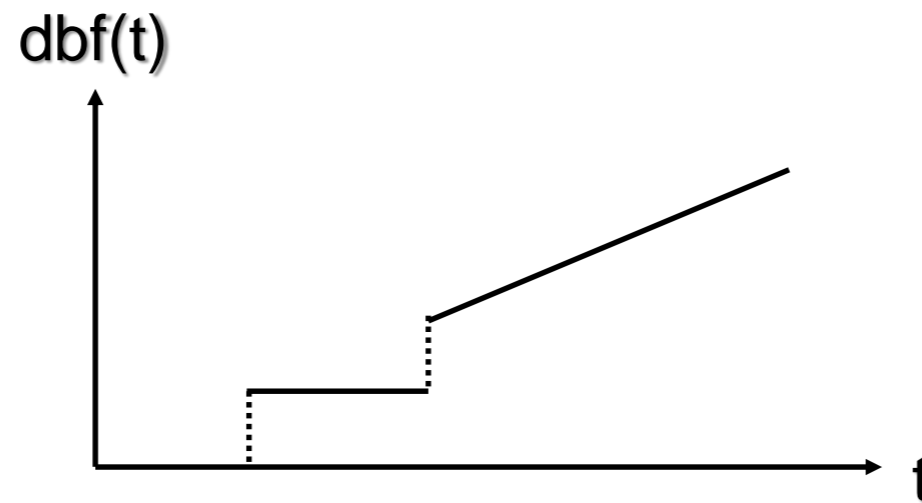
Based on demand bound functions



Exact test, **Pseudo-polynomial**
complexity if $\sum \alpha_i < 1$

Issues with admission control

- Some approximations exist to limit the computational complexity of the admission-test
- They are based on approximating the demand-bound function with a fixed number of discontinuities (Fisher et al., 2006)



**Polynomial-time complexity
(sufficient test)**

Questions



How to modify the **replenishment rules** for obtaining a **better** maximum-service **delay**?



Is it possible to achieve a maximum service **delay** equal to $\Delta = D + T - 2Q$?



How to guarantee a **bounded-delay** portion in presence of **shared-resources**?



Which **admission control test** should be used for admitting reservations?

THE QUESTION



How to implement a new **Hard Constant Bandwidth Server** supporting **constrained-deadlines**?

SKETCH OF SOLUTION:
SHADOW BUDGETING

Sketch of solution

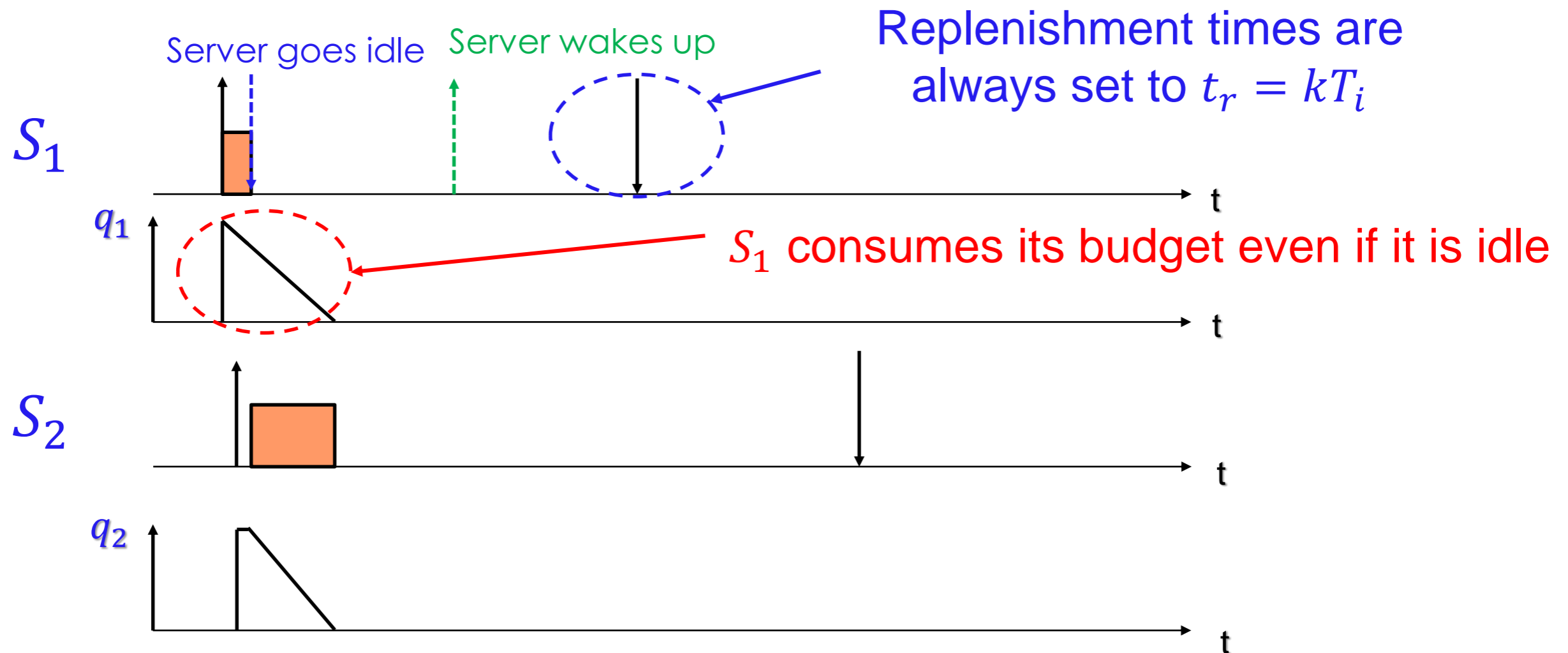
- The results proposed by Biondi et al. for **real-time self-suspending tasks** can be used to derive a solution

Alessandro Biondi, Alessio Balsini, and Mauro Marinoni,
“Resource reservation for real-time self-suspending tasks:
theory and practice” (RTNS 2015)

- According to their approach, whenever a server **should execute according to EDF** scheduling, it **consumes its budget** independently whether it is suspended or not

Shadow budget

- A similar approach can be adopted when a reservation goes idle:



Pro and Cons



Simplicity



Worst-case service delay is smallest as possible



Independent from the admission test



Lower throughput (average-case)



Still do not consider shared resources

What are we doing?

- Evaluation of different solutions
- Simulations
- Derive methodologies to increase the throughput
- Develop a solution to cope with shared resources
- Implement the new resource reservation server in Linux (SCHED_DEADLINE)

Thank you!

Daniel Casini
daniel.casini@sssup.it