



$\partial_t \psi + \frac{M}{\epsilon} \int_{\Omega} \frac{|u(x,t)|^2}{2} dx \Delta \psi + \int_{\Omega} \nabla \psi \cdot \nabla \psi = 0$ ,  $\int_{\Omega} \frac{|u(x,t)|^2}{2} dx = \int_{\Omega} \frac{|u(x,0)|^2}{2} dx = \rho_0(x)$ ,  $\psi(x,t) = e^{-i \int_{\Omega} \frac{|u(x,t)|^2}{2} dx} \psi_0(x)$

# Markov Chain Modelling of Probabilistic Real-Time Systems

Jasdeep Singh, Luca Santinelli, Guillaume Infantes, David Doose, Julien Brunel



retour sur innovation

# Real-Time System and Safety

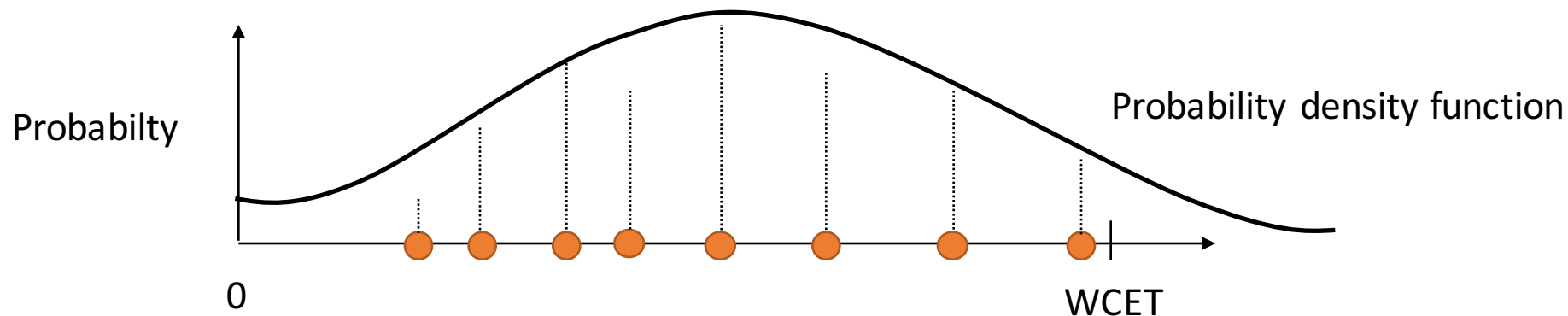
- Safety: ensure the correct working of R/T system; functional and non-functional aspects verified
- Modelling perspective: Ensure Task Deadlines are met
- R/T analysis required to verify this, in turn to ensure safety

# Why Probabilities

- Task execution time: WCET, worst case
- Not all instances are worst case
- Over pessimistic analysis
- Each execution time  $(0,WCET]$  has a probability of occurrence
- Representation of many possible execution times by a random variable
- Probability can quantify the pessimism

# Probabilistic Real-Time System (pRTS)

- R/T system with at least one probabilistic parameter
- WCET replaced by probabilistic WCET (pWCET)
- Each task has pWCET which is an upper bound probabilistic distribution of all possible execution times
- pWCET gives the probability that an execution  $\in (0, WCET]$  occurs



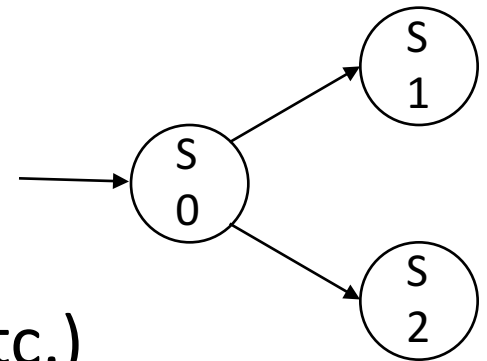
# Objective

- To perform probabilistic Schedulability analysis of pRTS
- Task set is given with one or more probabilistic parameters (pWCET, Deadline, Period)
- Model stochastic occurrences and interactions (execution, delays, preemptions, etc.)
- Provide probabilistic guarantees
- Probabilistically answer the question of safety

[Diaz et al. 2002, Carnevali et al 2014, Maxim et al 2013, Manolache et al 2004]

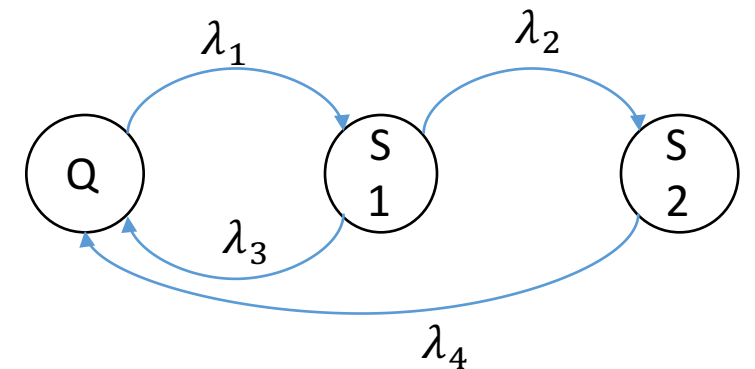
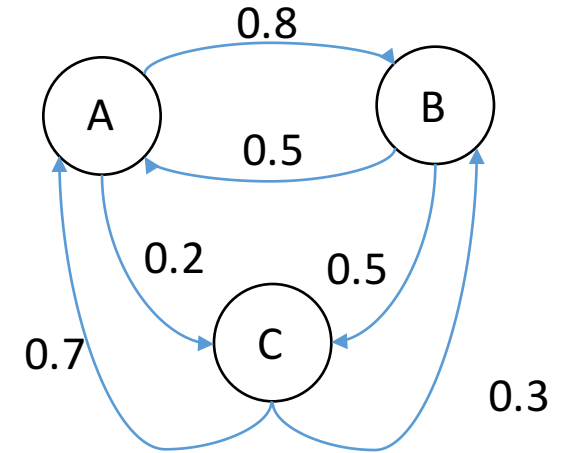
# Formal methods for pRTS

- Formal methods: mathematically proven techniques for specification and verification of a system
- Reliable: safe and accurate because of mathematical proofs available
- Model checking can be readily performed
- Model non-determinism and probability
- Handling continuous distribution (exponential, erlang, etc.)
- Feasible complexity



# A model approach

- Markov Chain (MC): set of states and transitions, memoryless property
- Continuous Time MC: transitions labelled with rates, duration of transition is exponentially distributed (choice and probability)
- Other stochastic formal methods
  - Stochastic Petri Net: Problem too complex
  - Stochastic Timed Automata: inability to model probabilistic duration of execution

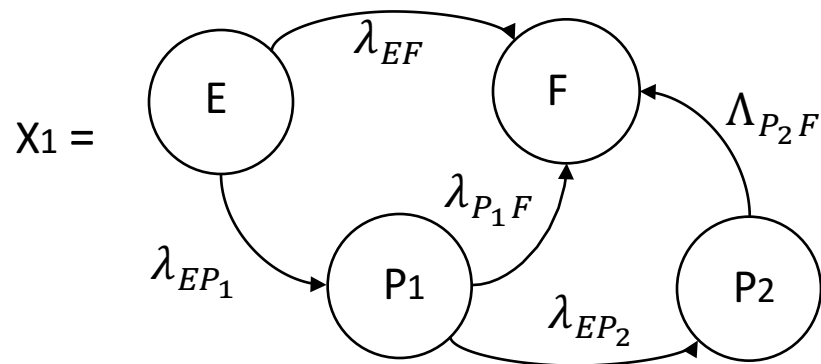
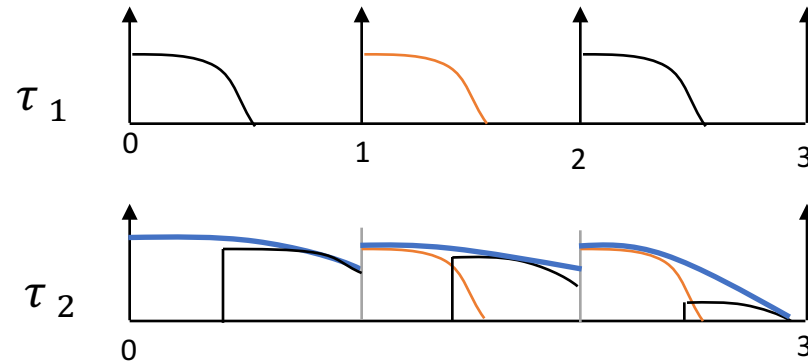


# Current Status

- Task set at input, tasks can be probabilistic and/or deterministic; execution time, period and deadline
- Probabilistic task have pWCET for execution time
- Scheduling policy assigns each job a priority (EDF, FP)
- Number of preemptions are known, since dealing with distributions with base  $(0, \infty)$



For  $\tau_2$ :



$$\Pr(\text{time} = 3; \text{state} = P_2) = \Pr(\text{DM})$$

- Execution times and delays encoded in transition rates
- Deadlines checked as property

# Limitations

- CTMC transitions are exponentially distributed
- Rigid time constraints like killing of job at deadline is not possible yet
- Convolutions are too heavy
- Hyperperiod delays not accounted for

# Open Problems and Future Work

- Current methods is in the process of publishing
- Next steps: Imposing rigid timing constraints
- Avoid convolution by different interpretation of the requirement
- Integration with random variables as limits
- Hybrid stochastic model (Automata+ CTMC; CTMC+DTMC, CTMC+CTMC): to impose hard timing constraints (preemption, killing job, etc.)

Thank you  
Questions